



Institut Mines-Télécom

# 6LoWPAN: Fragmentation & Reassembly, Frame Delivery Modes, & Fragment Forwarding

Georgios Z. PAPADOPOULOS, Associate Professor, IMT Atlantique

[georgios.papadopoulos@imt-atlantique.fr](mailto:georgios.papadopoulos@imt-atlantique.fr)

[www.georgiospapadopoulos.com](http://www.georgiospapadopoulos.com)

[www.youtube.com/c/gzpapadopoulos](https://www.youtube.com/c/gzpapadopoulos)

# SOMMAIRE

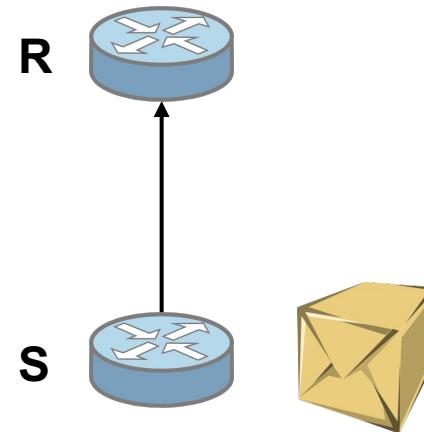
- 1. RFC 4944: 6LoWPAN (Per-Hop) Fragmentation & Reassembly**
- 2. RFC 4944: 6LoWPAN Frame Delivery Modes**
  - 2.1 Mesh-Under
  - 2.2 Route-Over
- 3. Route Over (Per-hop Fragmentation & Reassembly): Issues**
- 4. RFC 8930: 6LoWPAN Fragment Forwarding (6LFF)**

# Chapter 1

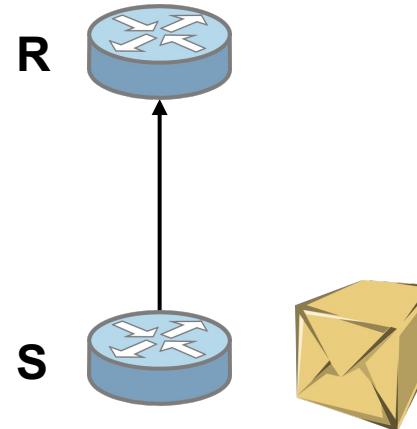
# RFC 4944: 6LoWPAN

# Fragmentation & Reassembly

RFC 4944



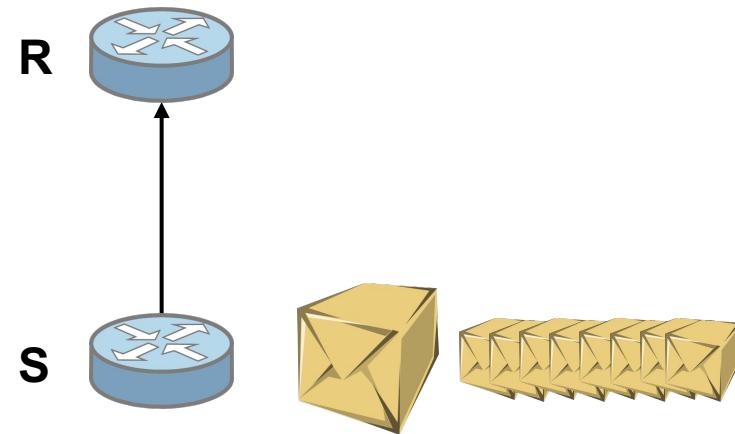
RFC 4944



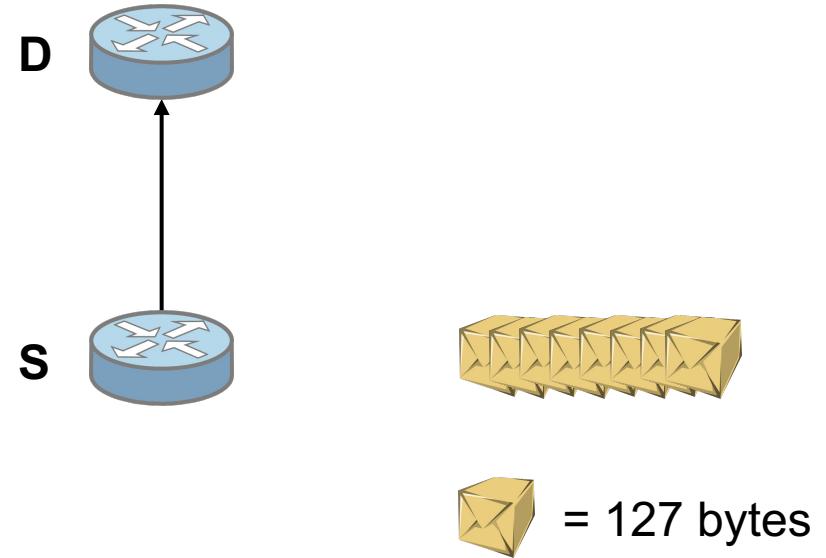
*When an IPv6 packet is larger than IEEE 802.15.4 MTU (127 bytes)*



RFC 4944

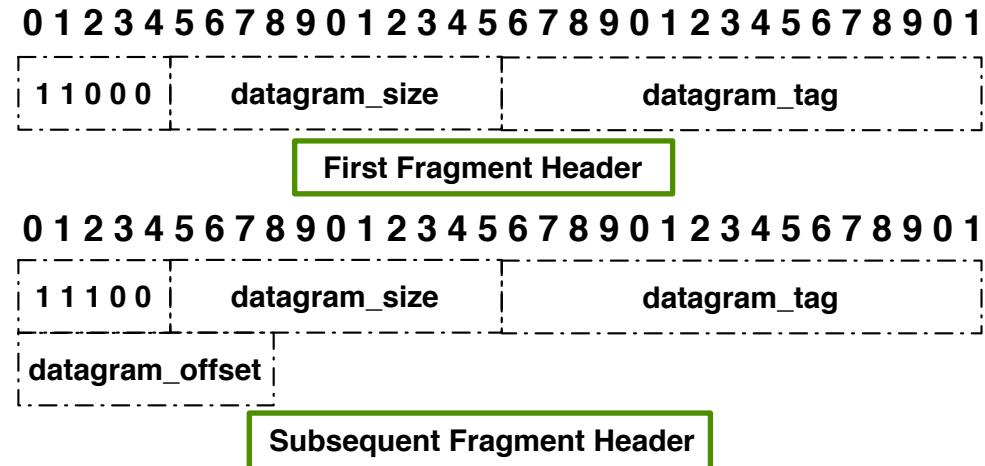


RFC 4944



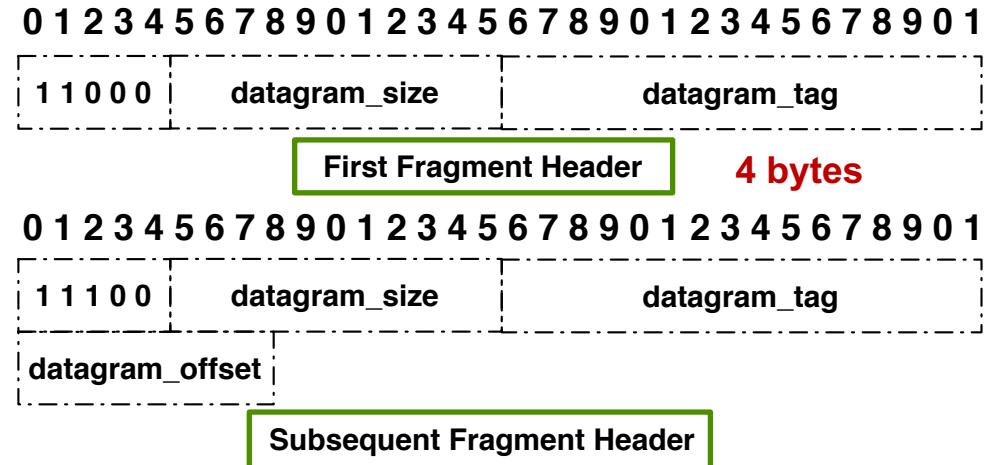
## Fragmentation Operation

RFC 4944



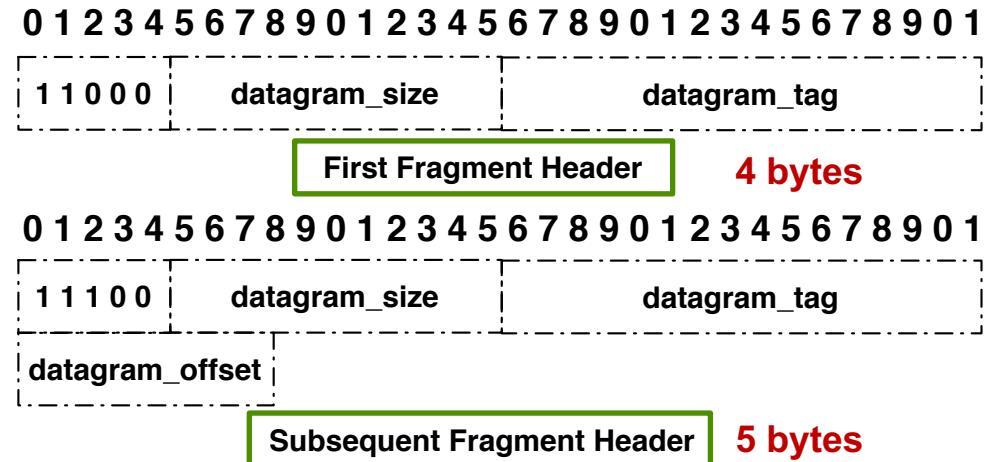
## Fragmentation Operation

RFC 4944



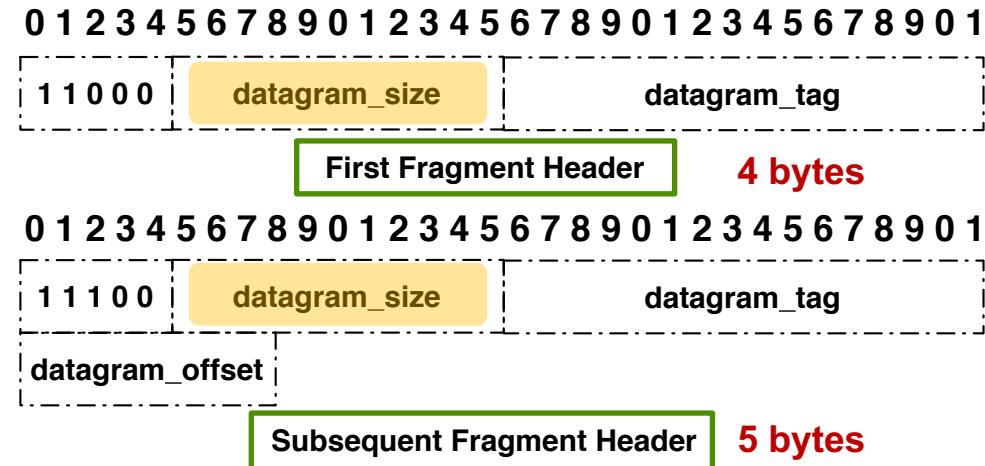
## Fragmentation Operation

RFC 4944



## Fragmentation Operation

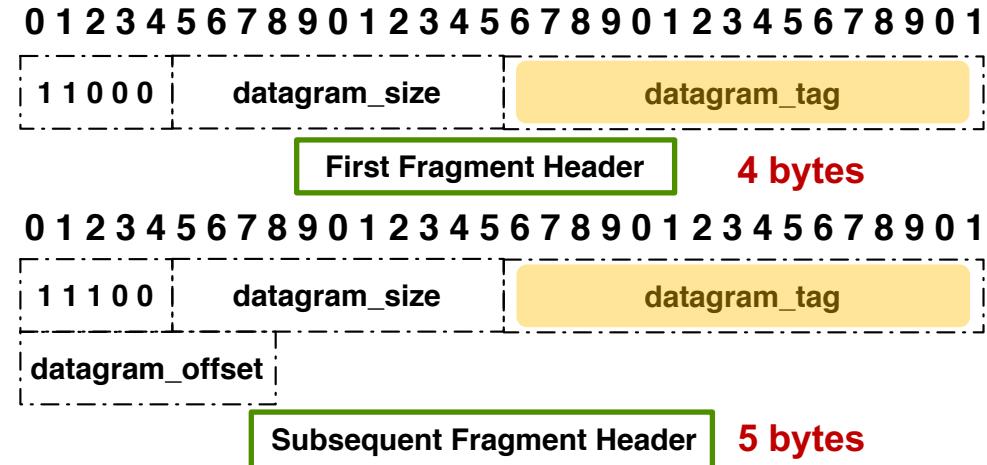
RFC 4944



1. The **datagram\_size** to identify the size of the IPv6 datagram.

## Fragmentation Operation

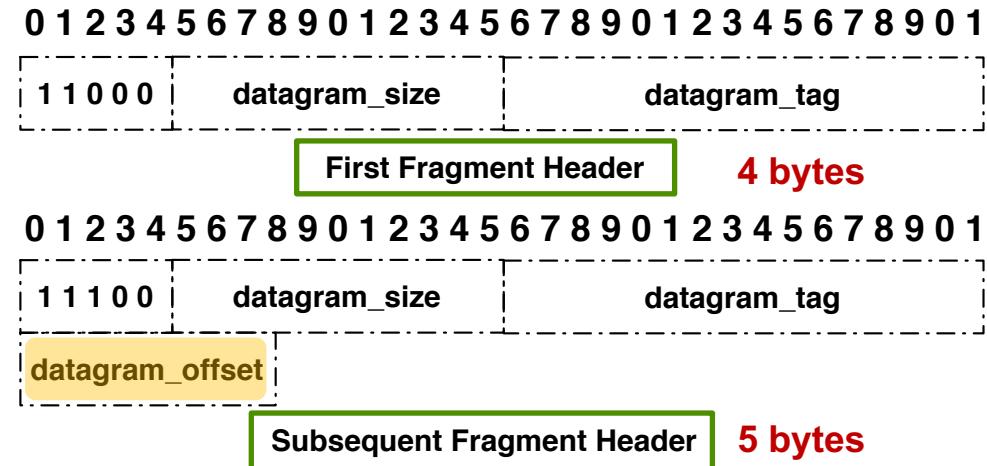
RFC 4944



2. The **datagram\_tag** to identify all fragments of a single datagram.

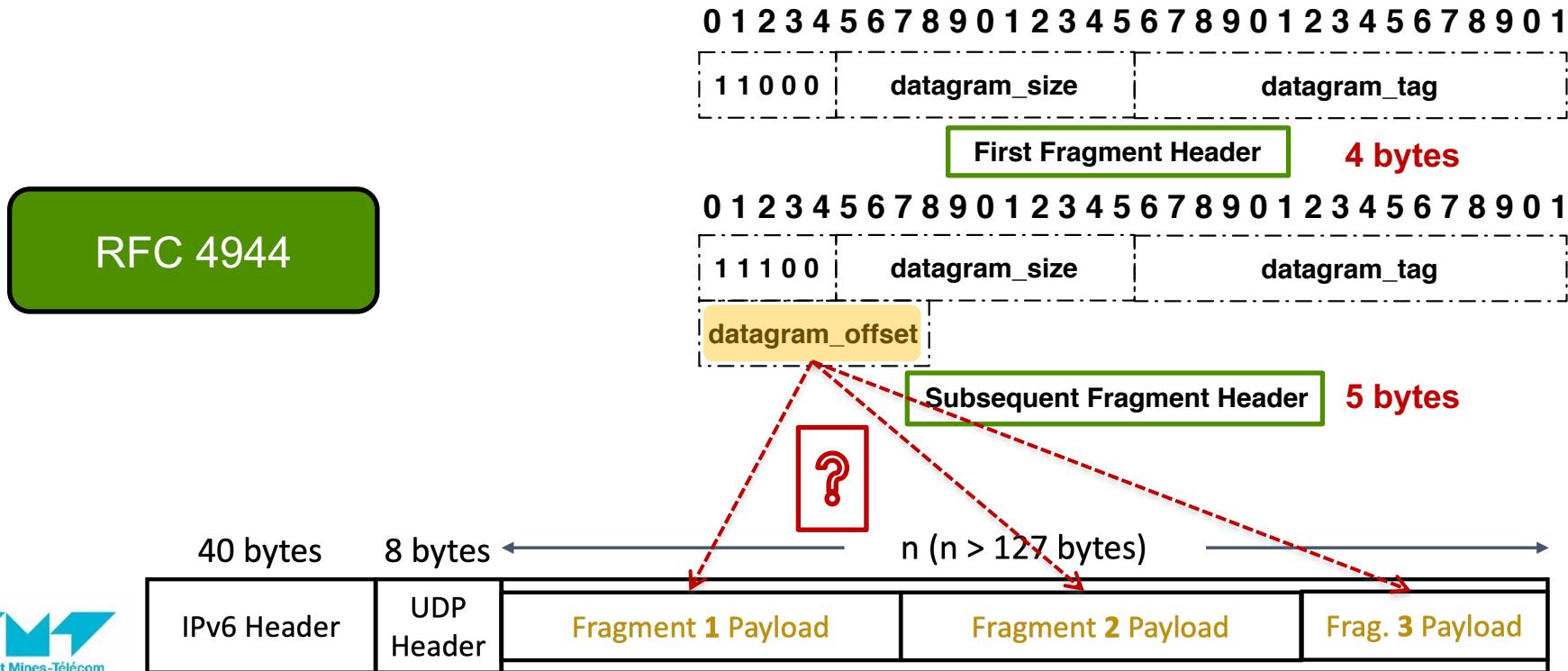
## Fragmentation Operation

RFC 4944



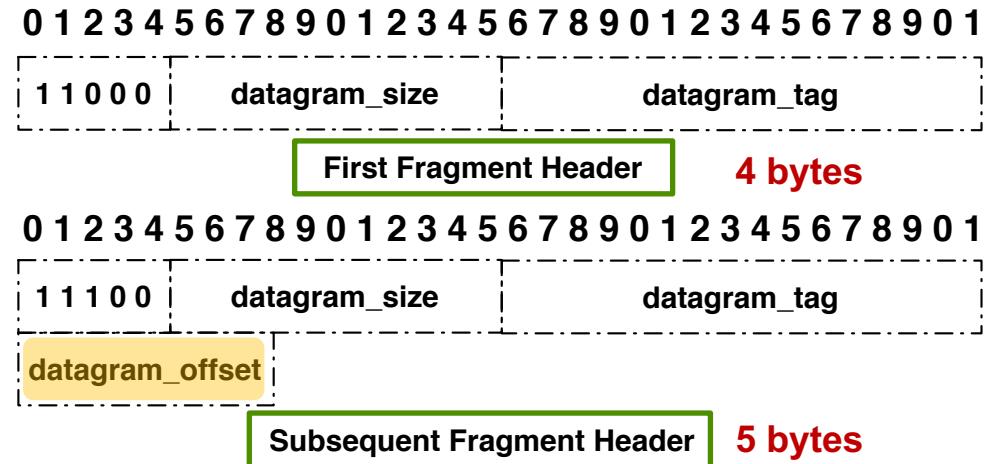
3. The **datagram\_offset** to identify the location of the received fragment.

## Fragmentation Operation

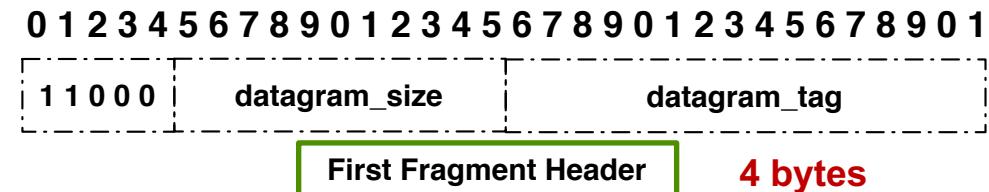


## Fragmentation Operation

RFC 4944

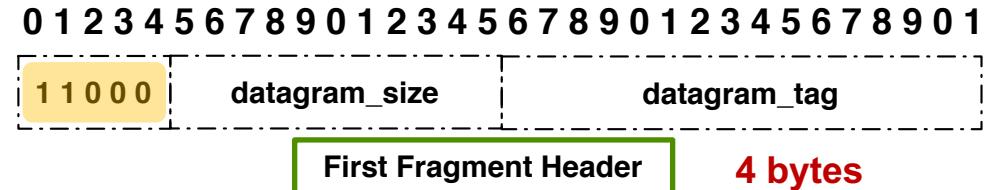


## Fragmentation Operation



RFC 4944

## Fragmentation Operation

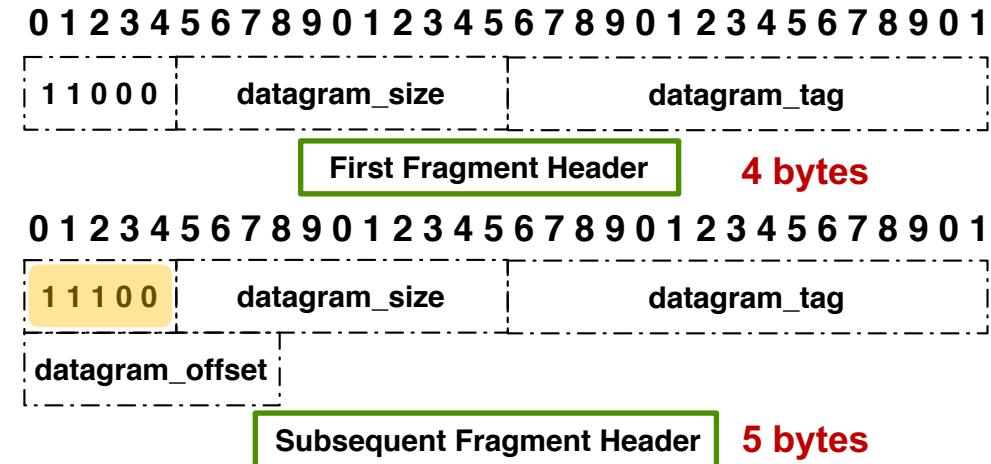


RFC 4944

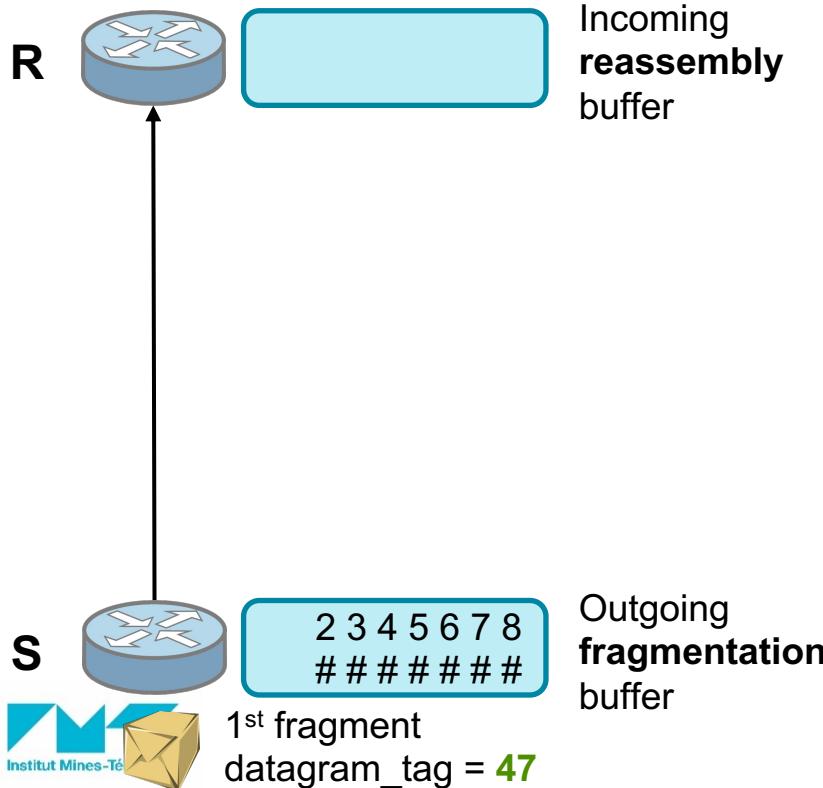
*Dispatch Value Bit Pattern → the Fragmentation Type*

## Fragmentation Operation

RFC 4944

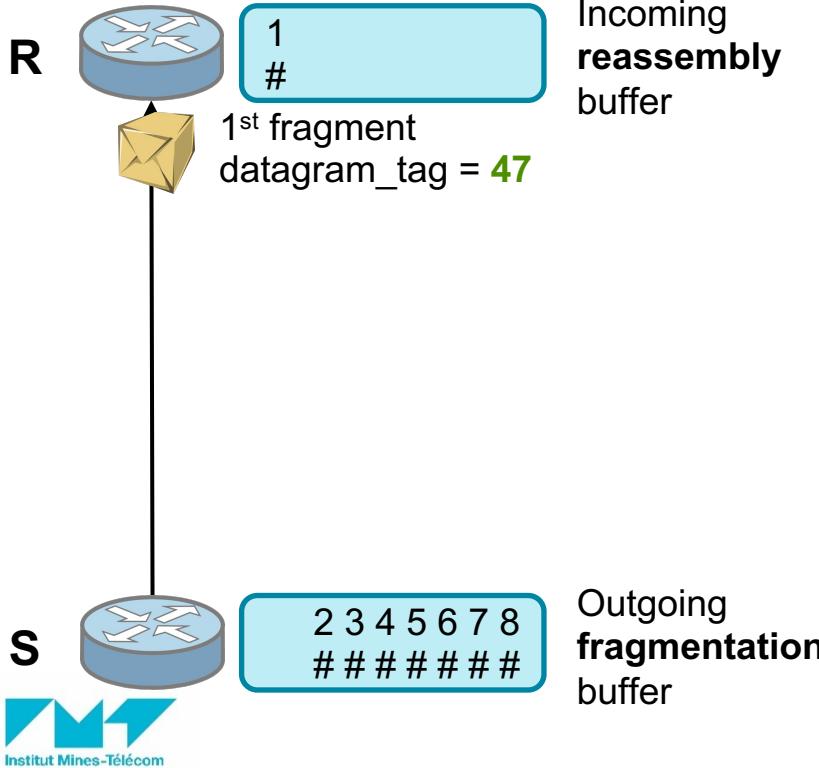


## Reassembly Operation



**The receiving node R:**

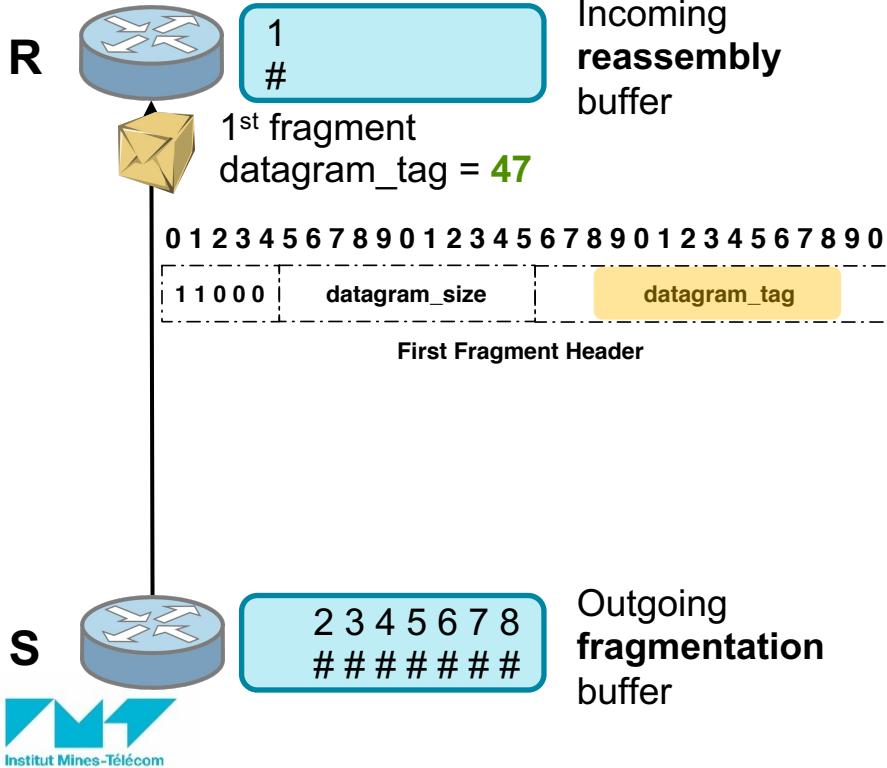
## Reassembly Operation



### The receiving node **R**:

- ▶ Initiates the reassembly operation to reconstruct the original IPv6 datagram.

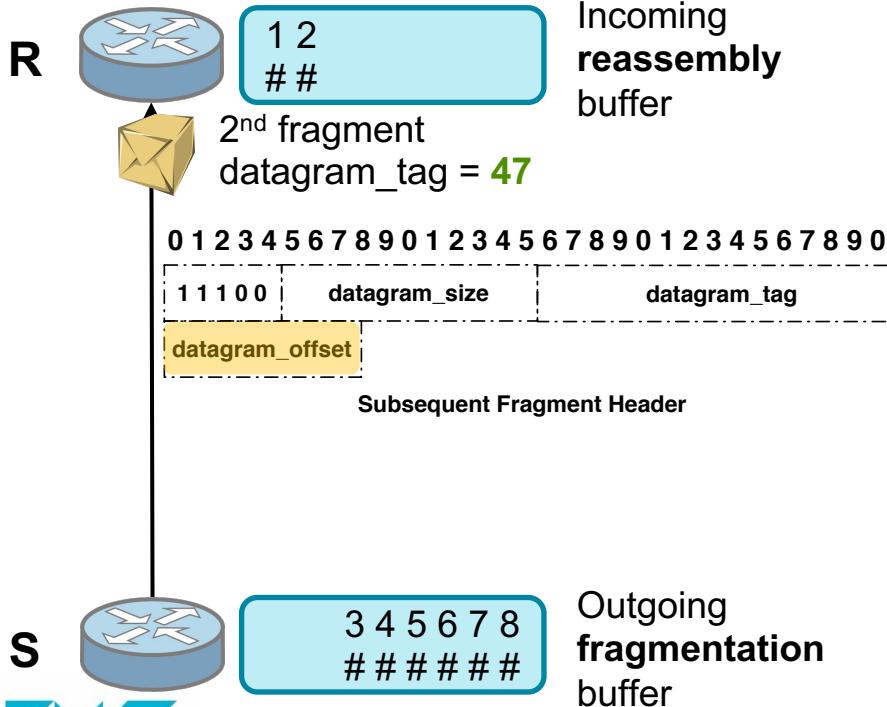
## Reassembly Operation



### The receiving node **R**:

1. Checks the **datagram\_tag** to identify the fragments that belong to a given IPv6 datagram.

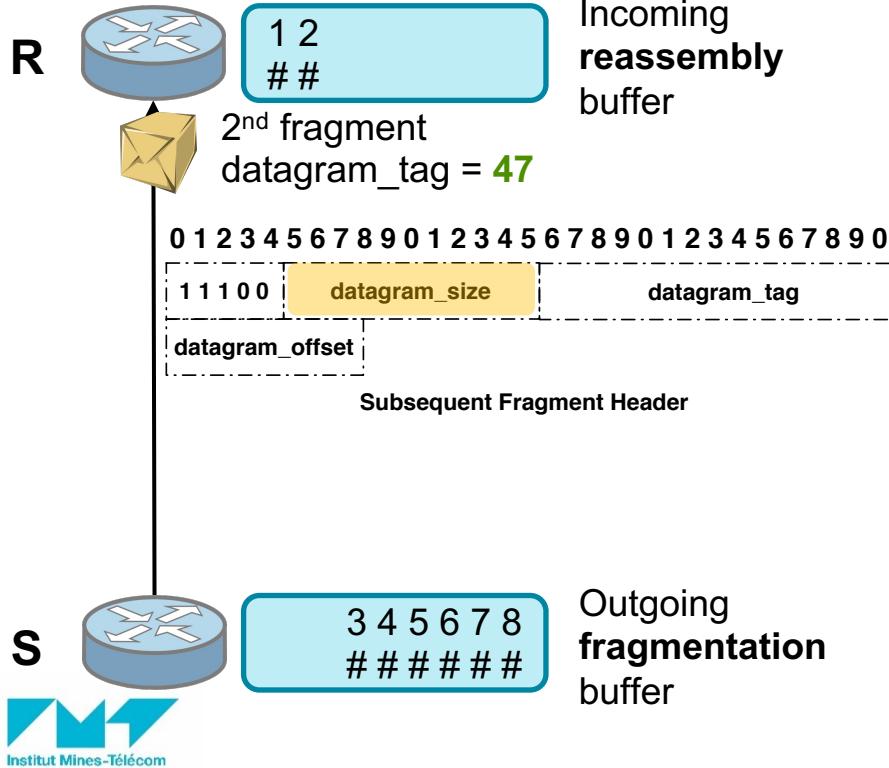
## Reassembly Operation



### The receiving node R:

2. Checks the ***datagram\_offset*** to determine the location of the received individual fragment.

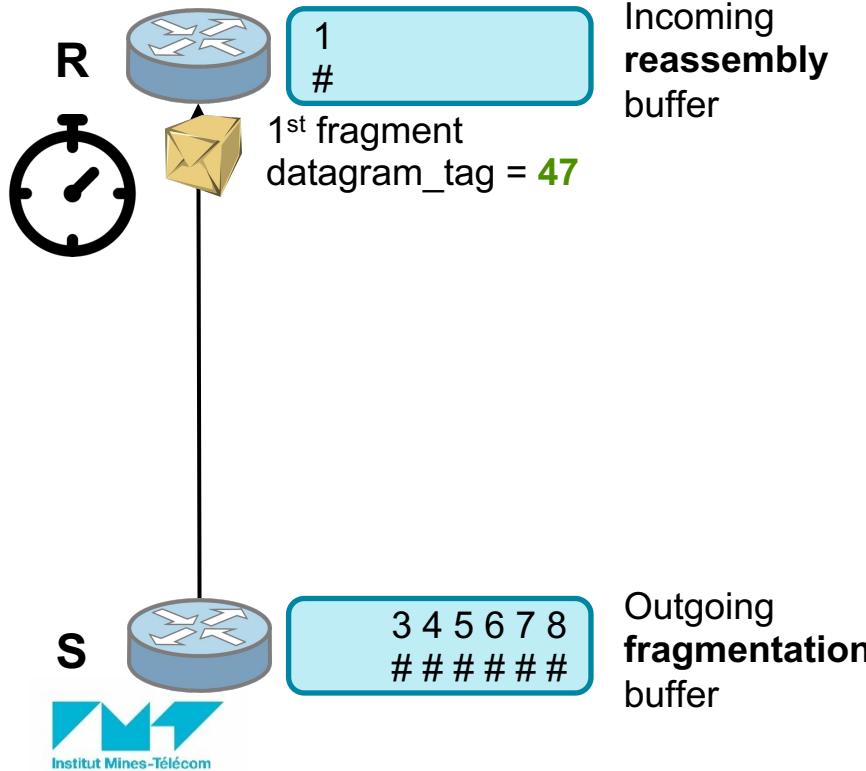
## Reassembly Operation



### The receiving node **R**:

3. Checks the **datagram\_size** to identify the size of the original unfragmented datagram.

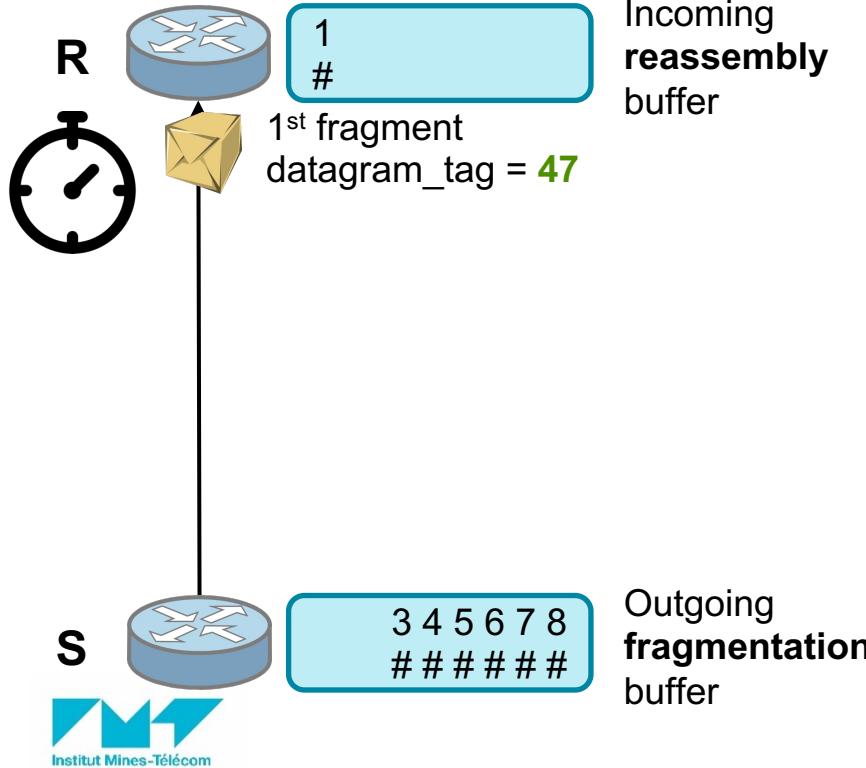
## The Reassembly Timer



**The receiving node R:**

- ▶ Sets a *reassembly timer*

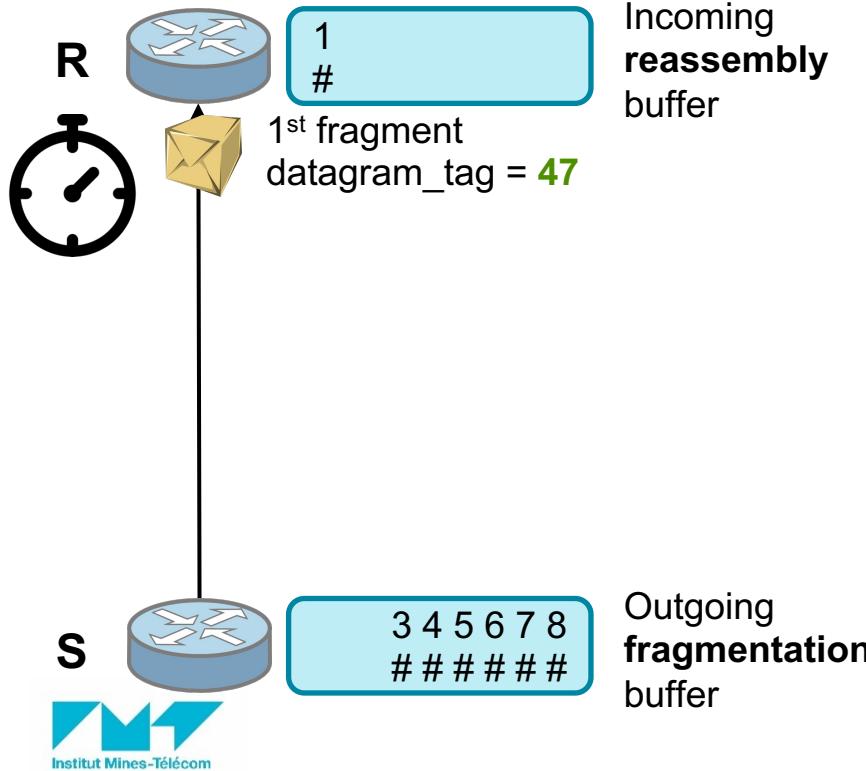
## The Reassembly Timer



## The receiving node **R**:

- ▶ Sets a reassembly timer, **which is configured to a maximum of 60 seconds**

## The Reassembly Timer



### The receiving node **R**:

- ▶ Sets a reassembly timer, which is configured to a maximum of 60 seconds **to allow discarding the partially reassembled packet after the timeout.**

# Chapter 2

# RFC 4944: 6LoWPAN

# Frame Delivery Modes

RFC 4944

## 6LoWPAN Frame Delivery modes:

- ▶ Mesh-Under.
- ▶ Router-Over.

RFC 4944

## 6LoWPAN Frame Delivery modes:

- ▶ Mesh-Under: Layer 2 based on MAC.
- ▶ Router-Over.

RFC 4944

## 6LoWPAN Frame Delivery modes:

- ▶ Mesh-Under: Layer 2 based on MAC.
- ▶ Router-Over: Layer 3 based on IP.

## Mesh-Under

### Mesh-Under:

- ▶ The ***routing*** and ***forwarding*** operations are performed at 6LoWPAN Adaptation layer.

RFC 4944

## Mesh-Under

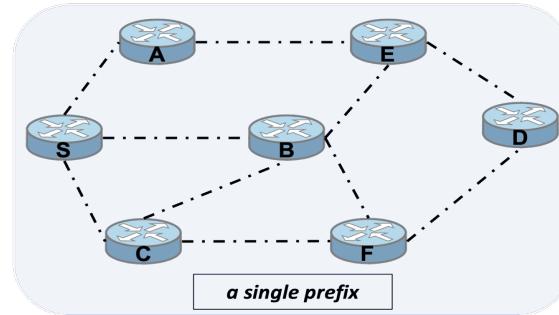
RFC 4944

### Mesh-Under:

- ▶ The ***routing*** and ***forwarding*** operations are performed at 6LoWPAN Adaptation layer.
- ▶ The IPv6 header ***does not*** need to be unpacked.

**Mesh-Under****RFC 4944****Mesh-Under:**

- ▶ The **routing** and **forwarding** operations are performed at 6LoWPAN Adaptation layer.
- ▶ The IPv6 header **does not** need to be unpacked.
- ▶ All nodes share the same prefix.



**Mesh-Under: Overview**

RFC 4944

To forward the received fragments,  
a node requires:

- ▶ The ***final destination address***, the link-layer address of the Final Destination.

## Mesh-Under: Overview

RFC 4944

To forward the received fragments,  
a node requires:

- ▶ The ***final destination address***, the link-layer address of the Final Destination.
- ▶ The ***originator address***, the link-layer address of the Originator.

**Mesh-Under: Overview**

RFC 4944

To forward the received fragments,  
a node requires:

- ▶ The ***final destination address***, the link-layer address of the Final Destination.
- ▶ The ***originator address***, the link-layer address of the Originator.

originator address, final address

**Mesh-Under: Mesh Addressing Type and Header****RFC 4944**

To forward the received fragments,  
a node requires:

- ▶ The ***final destination address***, the link-layer address of the Final Destination.
- ▶ The ***originator address***, the link-layer address of the Originator.

0	1	2	3	4	5	6	7	8	9	0	1	2	3	4	5	0	1	2	3
1	0									originator address, final address (16 or 64-bit)									

**Mesh-Under: Mesh Addressing Type and Header****RFC 4944**

To forward the received fragments, a node requires:

- ▶ The **final destination address**, the link-layer address of the Final Destination.
- ▶ The **originator address**, the link-layer address of the Originator.
- ▶ **Hops Left** (4 bits), a layer-2 equivalent of an IPv6 Hop Limit.

0	1	2	3	4	5	6	7	8	9	0	1	2	3	4	5	0	1	2	3
1	0																		

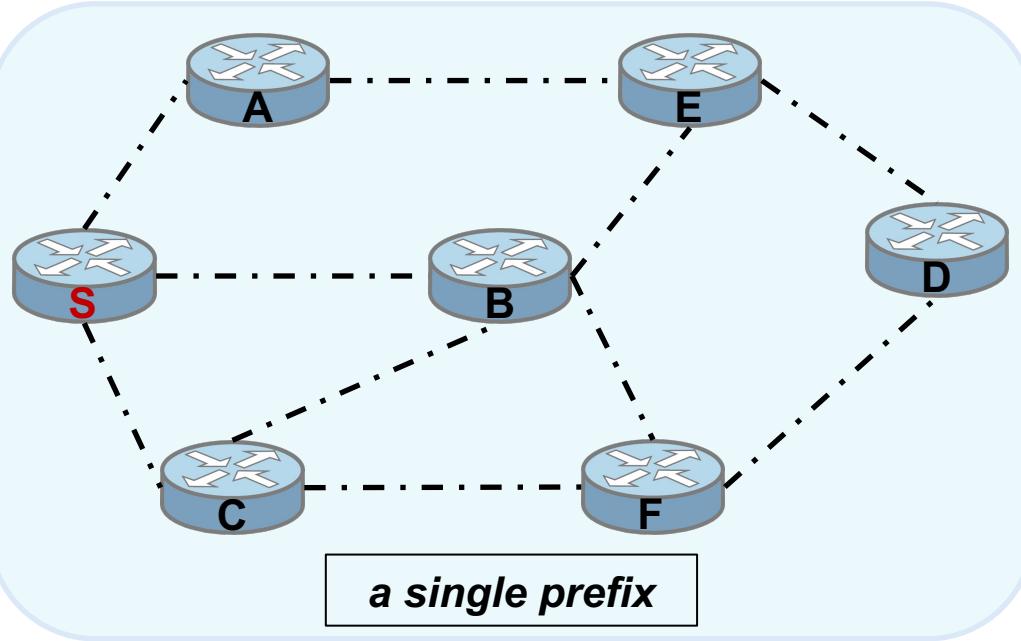
Hops Left      originator address, final address (16 or 64-bit)

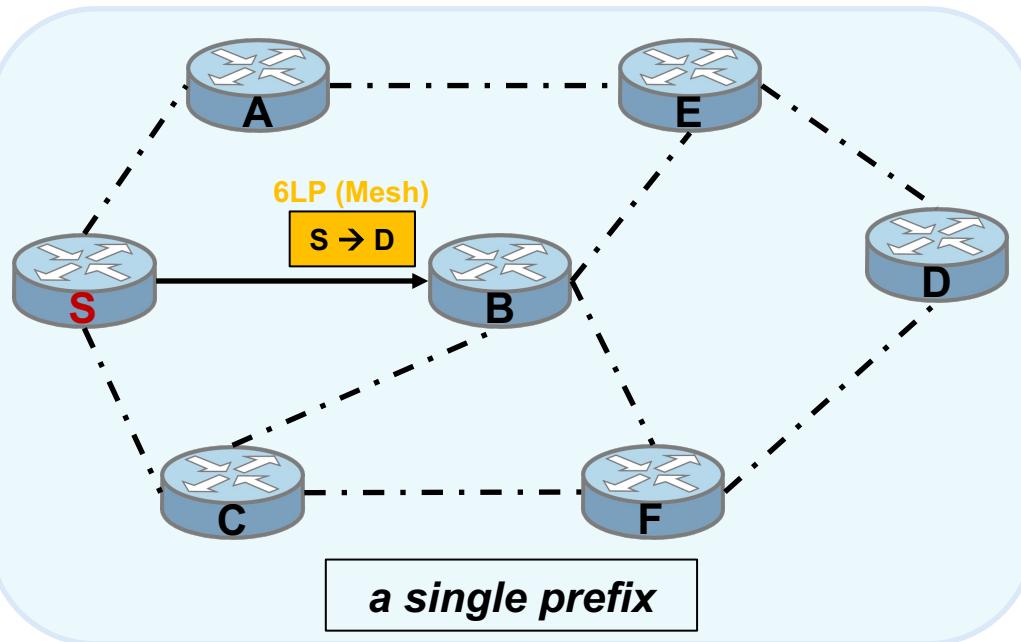
**Mesh-Under: Mesh Addressing Type and Header****RFC 4944**

**To forward the received fragments,  
a node requires:**

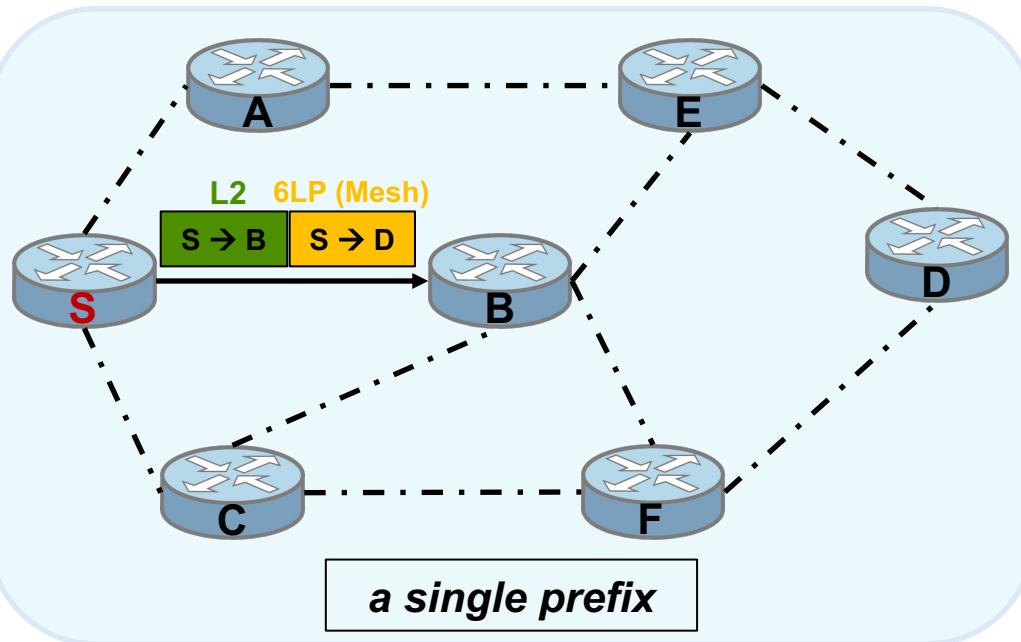
- ▶ The **final destination address**, the link-layer address of the Final Destination.
- ▶ The **originator address**, the link-layer address of the Originator.
- ▶ **Hops Left** (4 bits), a layer-2 equivalent of an IPv6 Hop Limit.
- ▶ **V** and **F bits** to indicate whether the addresses are 16-bit short or 64-bit EUI-64.

0	1	2	3	4	5	6	7	8	9	0	1	2	3	4	5	0	1	2	3
1	0	V	F		Hops Left														

**Mesh-Under: Operation Example**

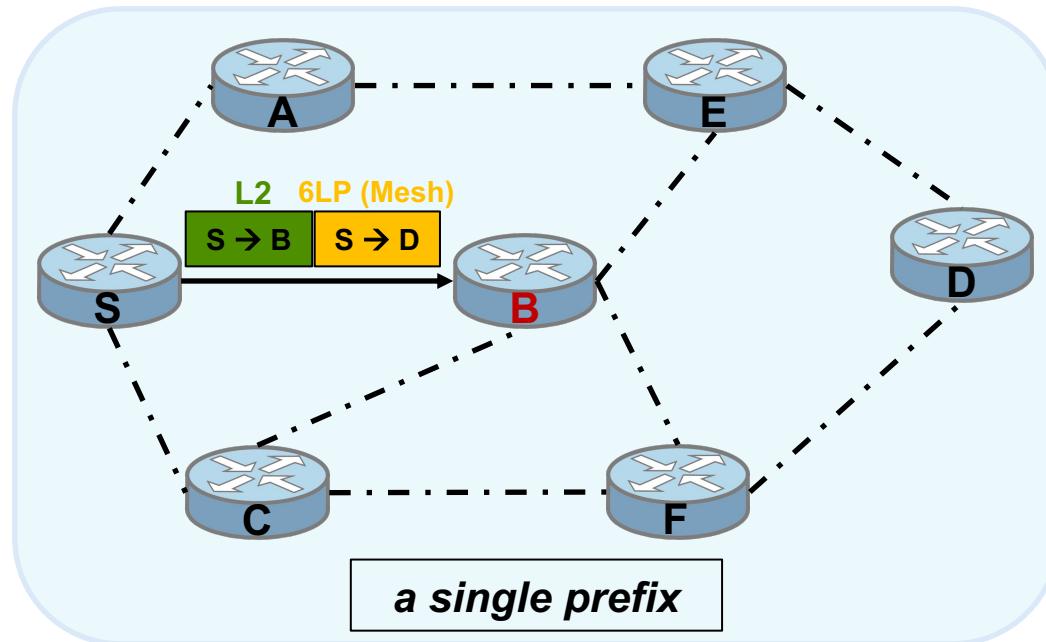
**Mesh-Under: Operation Example**

**6LP:** 6LoWPAN Adaptation Layer (i.e., Mesh Addressing Type & Header)

**Mesh-Under: Operation Example**

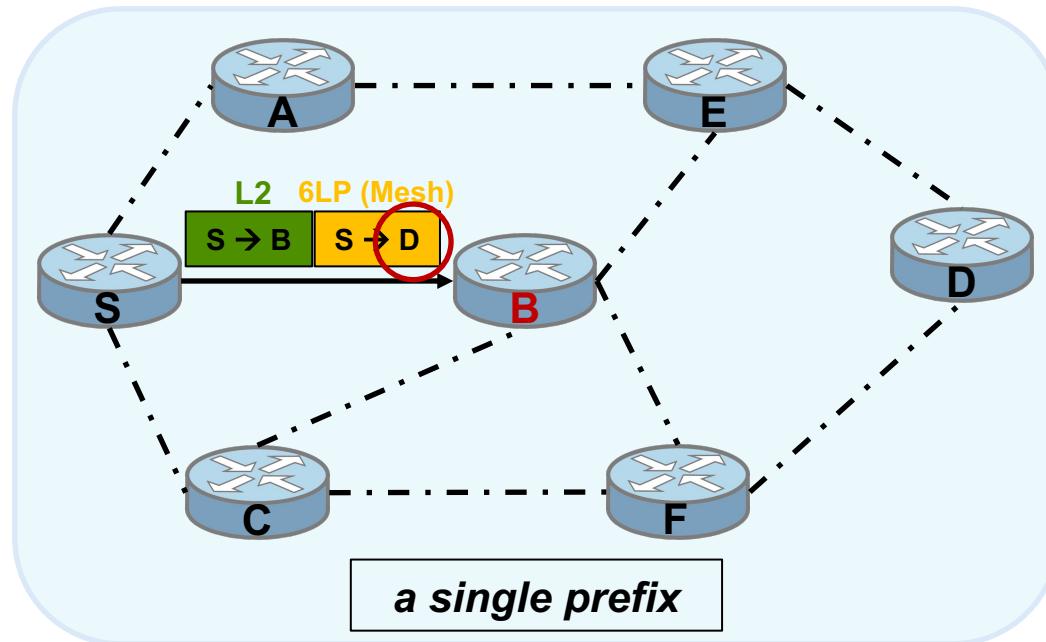
**6LP:** 6LoWPAN Adaptation Layer (i.e., Mesh Addressing Type & Header)

**L2:** Layer 2 (e.g., IEEE Std 802.15.4)

**Mesh-Under: Operation Example**

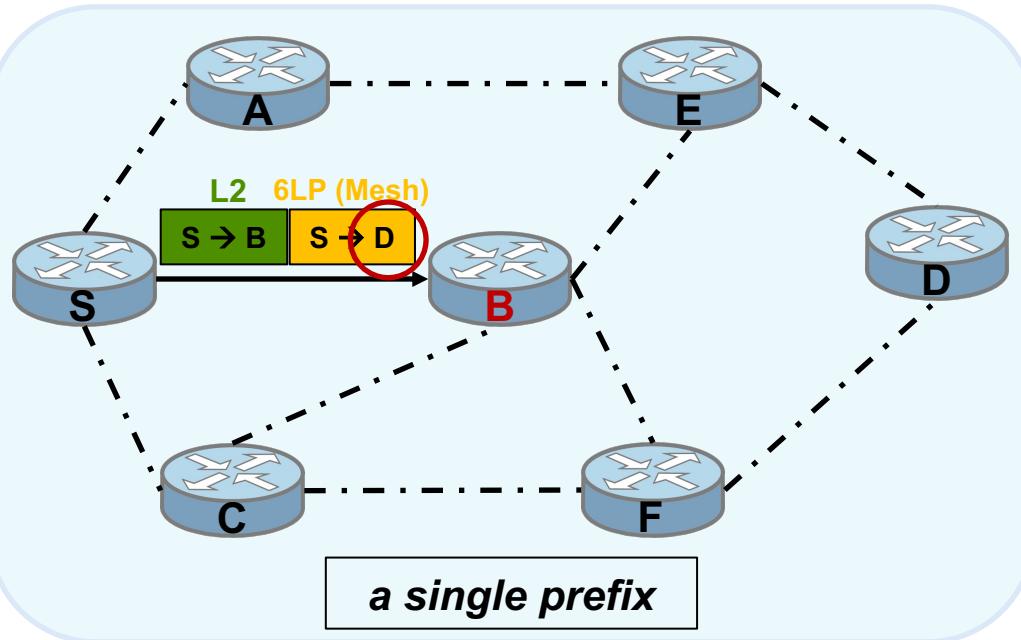
**6LP:** 6LoWPAN Adaptation Layer (i.e., Mesh Addressing Type & Header)

**L2:** Layer 2 (e.g., IEEE Std 802.15.4)

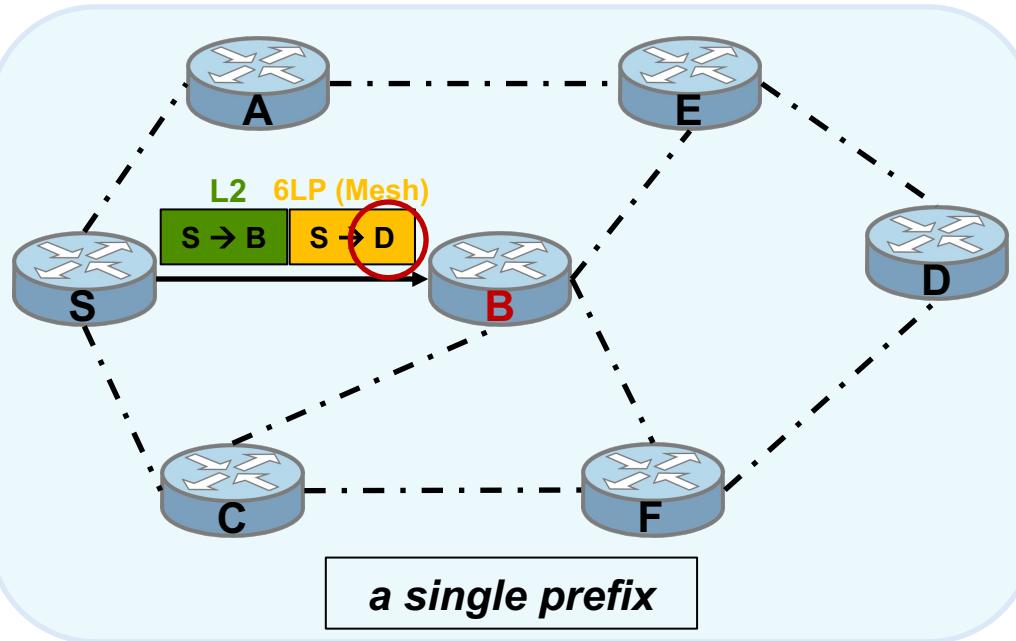
**Mesh-Under: Operation Example**

**6LP:** 6LoWPAN Adaptation Layer (i.e., Mesh Addressing Type & Header)

**L2:** Layer 2 (e.g., IEEE Std 802.15.4)

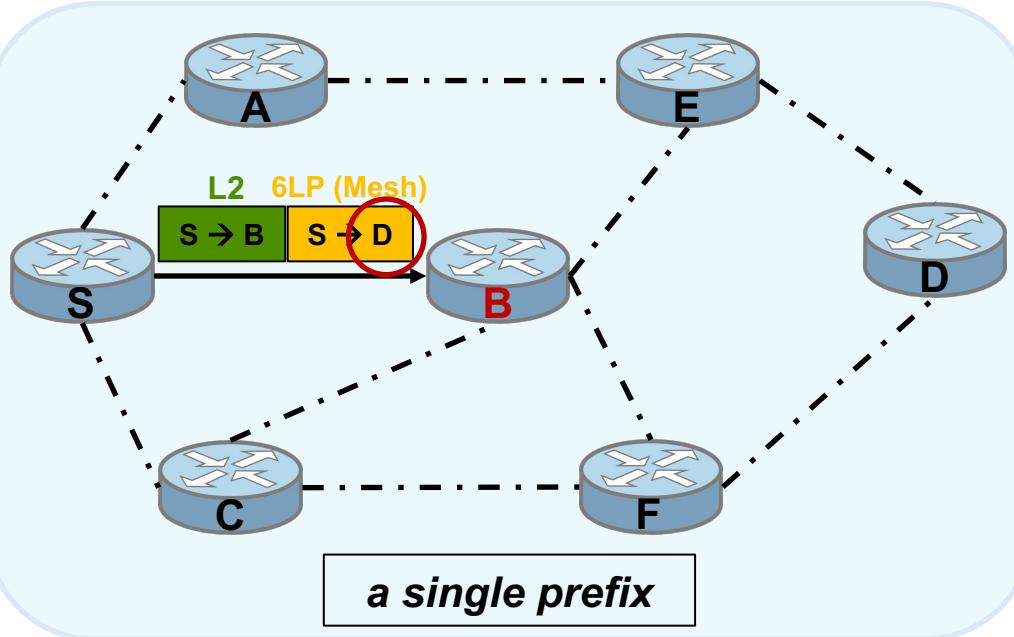
**Mesh-Under: Operation Example**

1. If the node **is** the final destination, → it initiates the reassembly of the IPv6 packet.

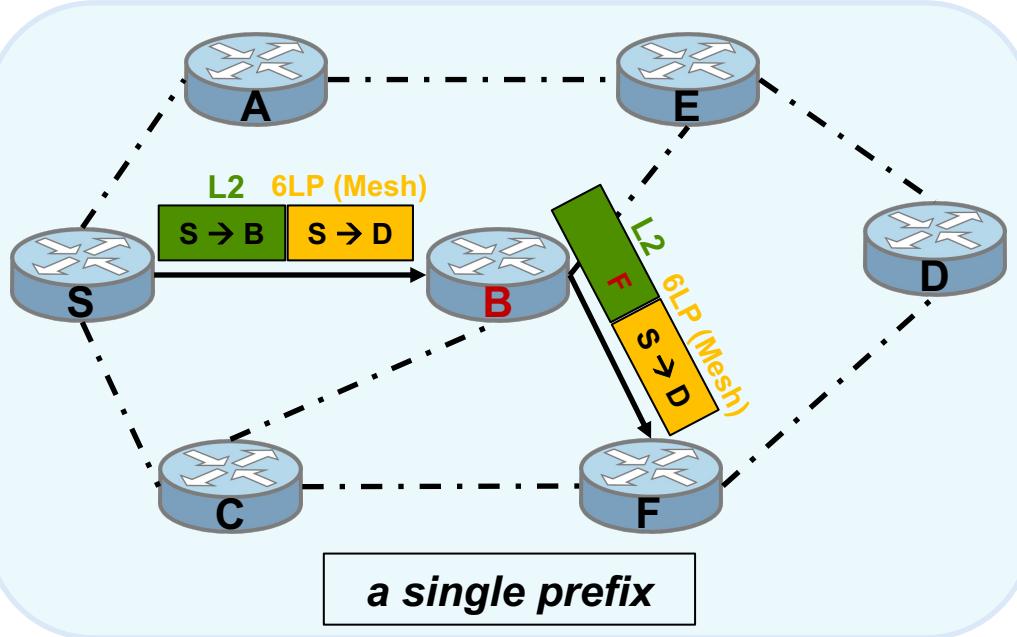
**Mesh-Under: Operation Example**

2. If it *is not*, → it reduces the "Hops Left",

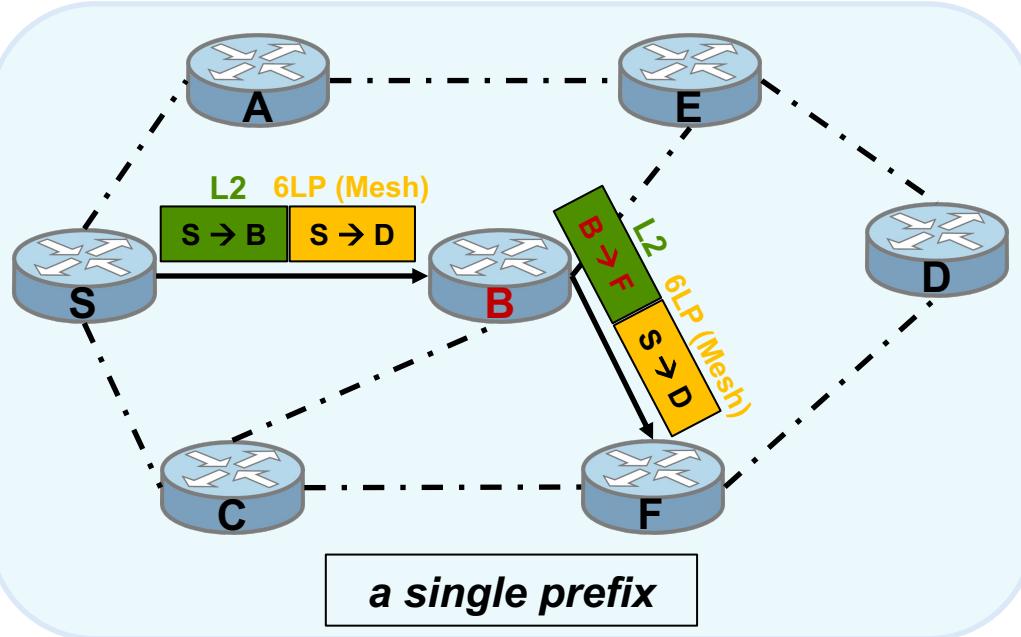
## Mesh-Under: Operation Example



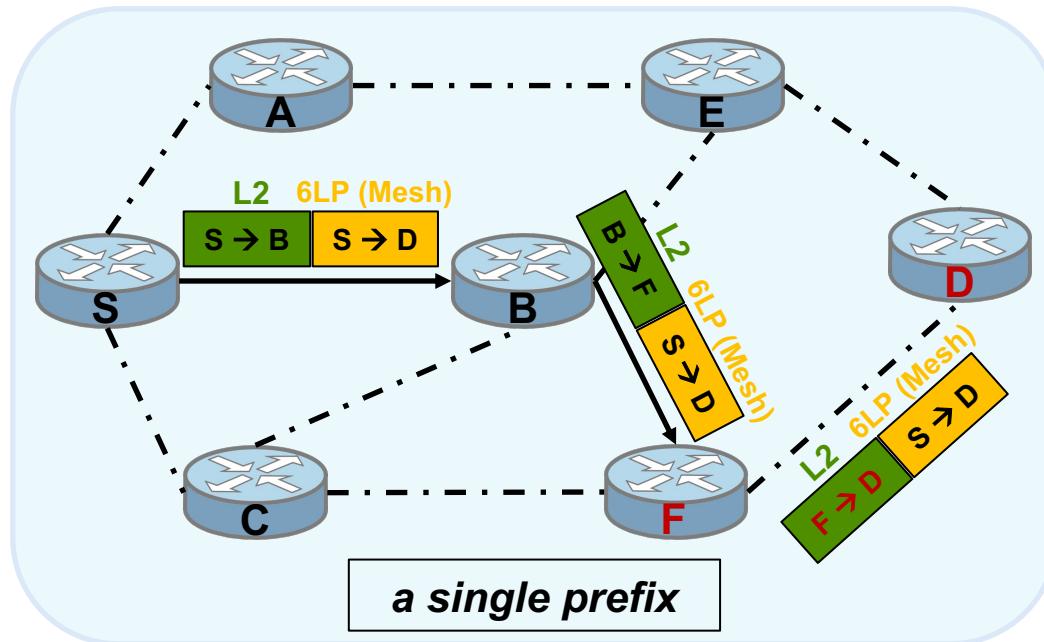
2. If it **is not**, →  
it reduces the "Hops Left",  
and if it **decremented to zero**, →  
it **discards** the frame.

**Mesh-Under: Operation Example**

3. Otherwise,
  - it sets F in the destination address field of the L2 header.

**Mesh-Under: Operation Example**

3. Otherwise,
  - it sets **F** in the destination address field of the **L2 header**.
  - Finally, it sets the **L2 source address to its own** and transmits the frame.

**Mesh-Under: Operation Example**

**6LP:** 6LoWPAN Adaptation Layer (i.e., Mesh Addressing Type & Header)

**L2:** Layer 2 (e.g., IEEE Std 802.15.4)

**Route-Over**

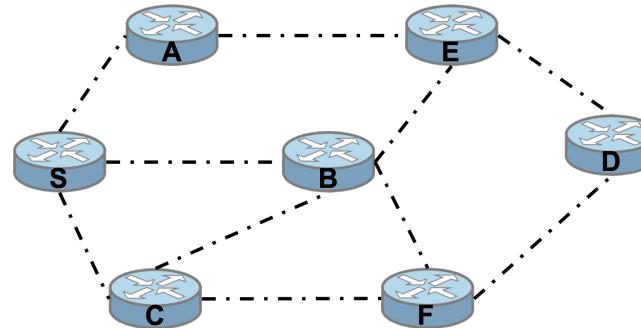
RFC 4944

**Route-Over:**

- ▶ The ***routing*** and ***forwarding*** operations are performed at Layer 3 based on IP.

**Route-Over****RFC 4944****Route-Over:**

- ▶ The ***routing*** and ***forwarding*** operations are performed at Layer 3 based on IP.
- ▶ Each link layer hop is an IP hop.



## Route-Over: Overview

Outgoing fragmentation buffer

**A**

# 4th fragment

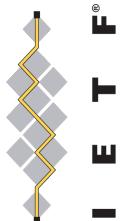


Incoming reassembly buffer

**B**

An IPv6 packet is expected to be:

- Reassembled at each intermediate node.

**C**

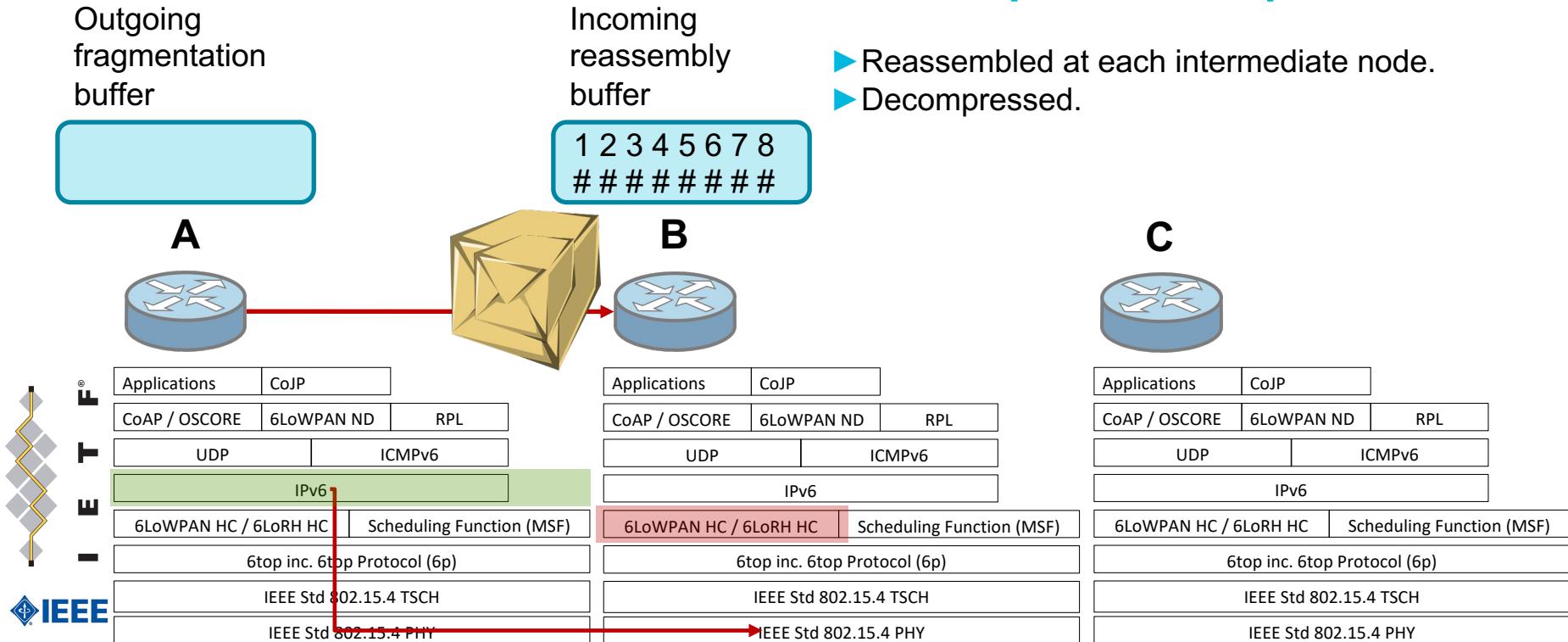
Applications	CoJP
CoAP / OSCORE	6LoWPAN ND
RPL	
UDP	ICMPv6
IPv6	
6LoWPAN HC / 6LoRH HC	Scheduling Function (MSF)
6top inc. 6top Protocol (6p)	
IEEE Std 802.15.4 TSCH	
IEEE Std 802.15.4 PHY	

Applications	CoJP
CoAP / OSCORE	6LoWPAN ND
RPL	
UDP	ICMPv6
IPv6	
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Applications	CoJP
CoAP / OSCORE	6LoWPAN ND
RPL	
UDP	ICMPv6
IPv6	
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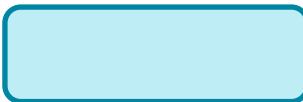


## Route-Over: Overview



## Route-Over: Overview

Outgoing fragmentation buffer



A



Incoming reassembly buffer



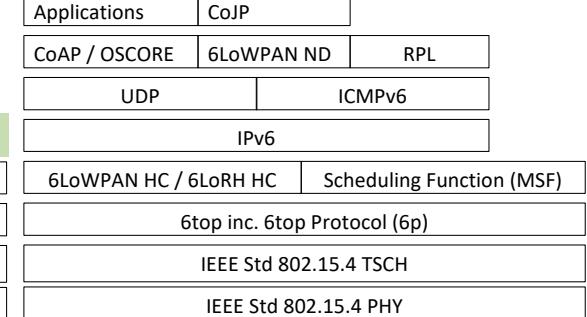
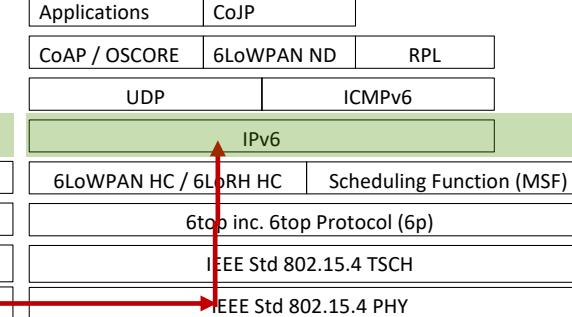
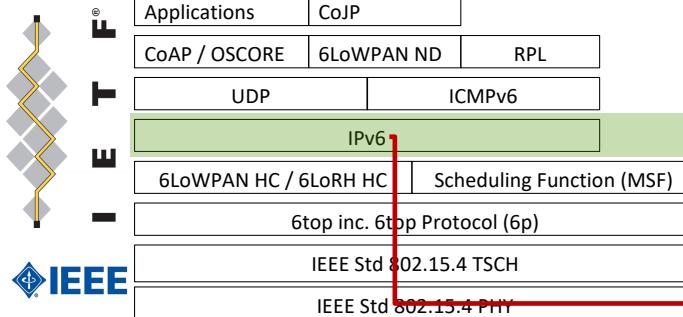
B



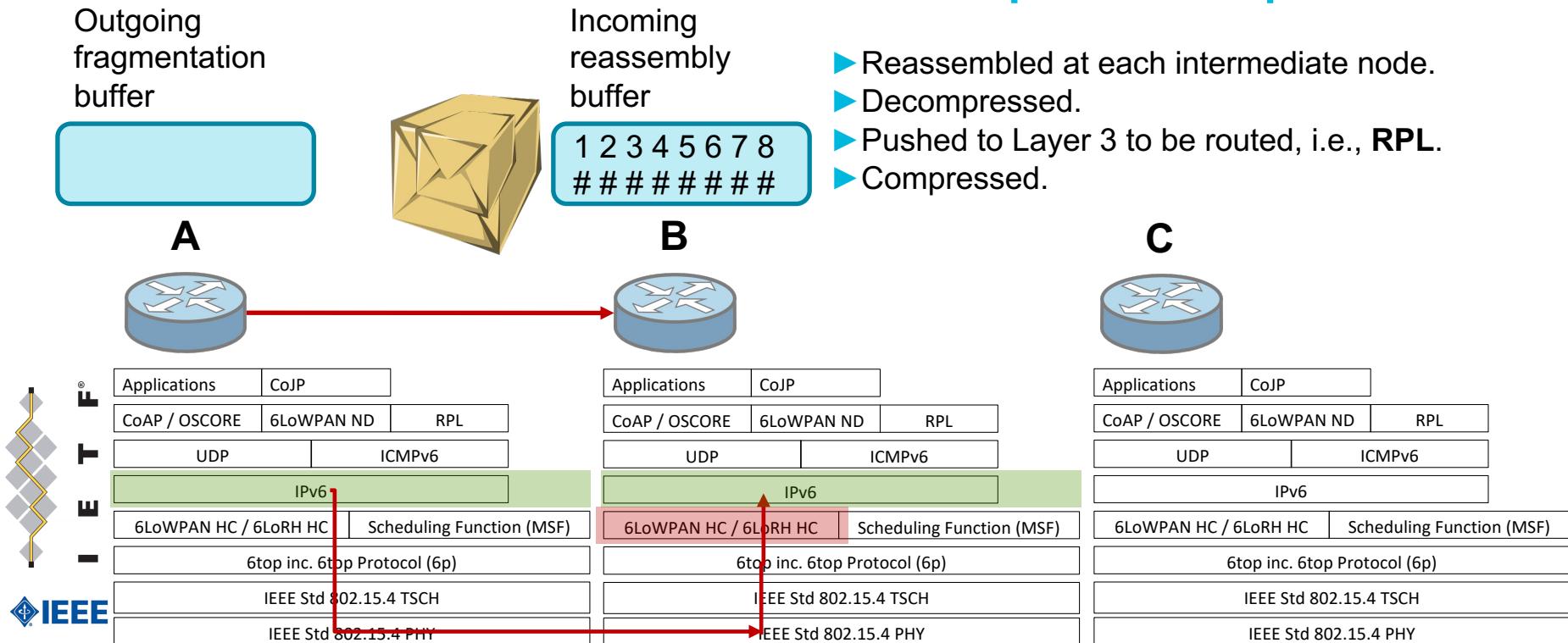
An IPv6 packet is expected to be:

- ▶ Reassembled at each intermediate node.
- ▶ Decompressed.
- ▶ Pushed to Layer 3 to be routed, i.e., **RPL**.

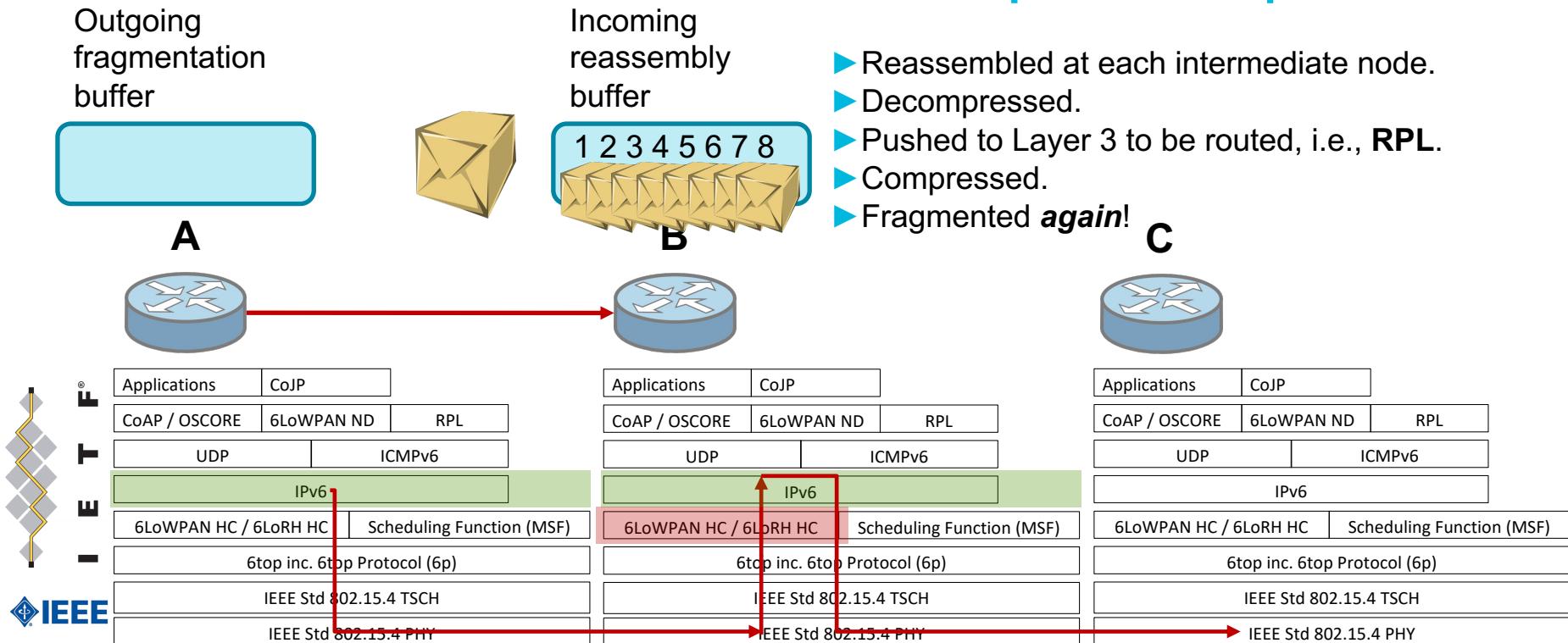
C

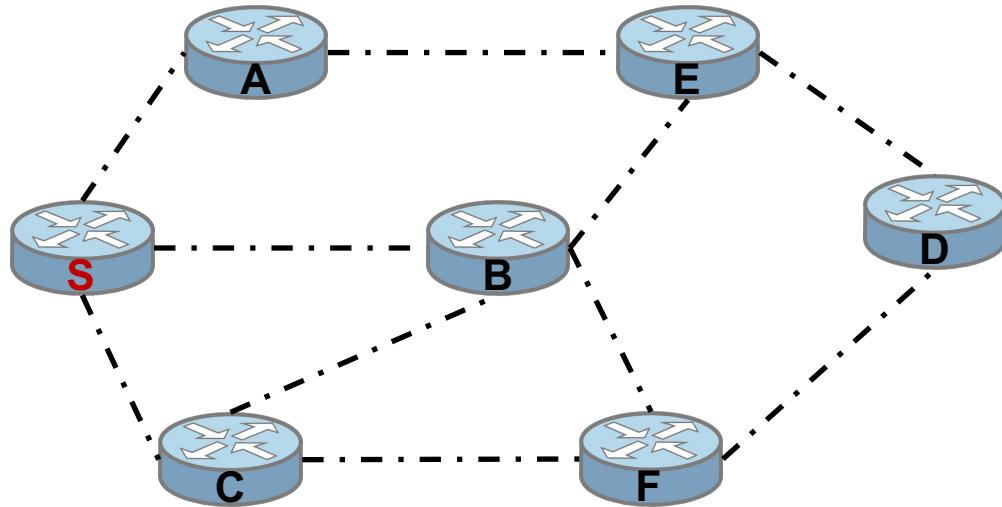


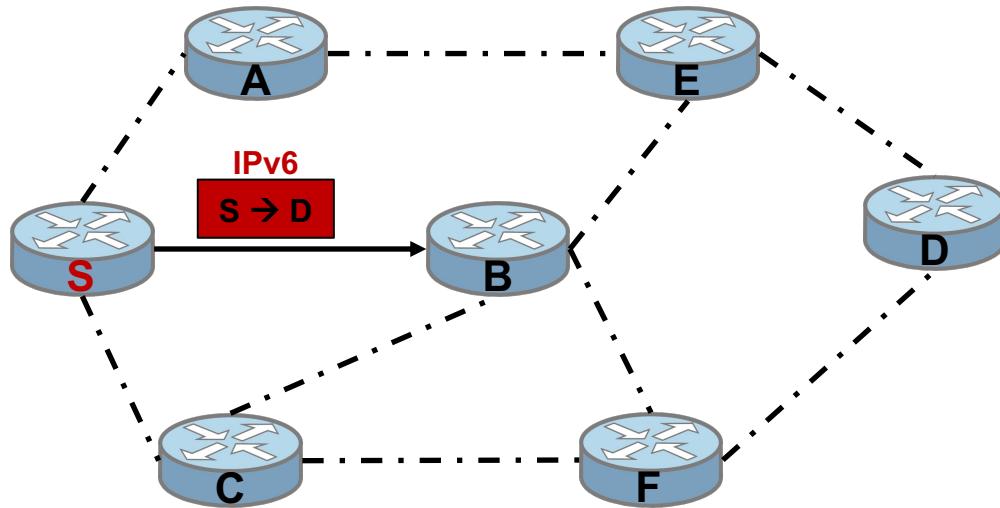
## Route-Over: Overview



## Route-Over: Overview

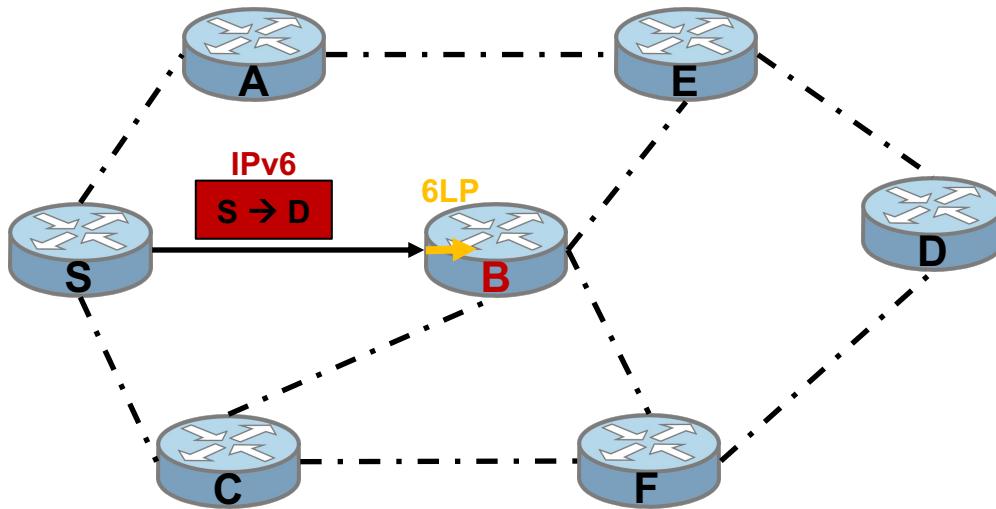


**Route-Over: Operation Example**

**Route-Over: Operation Example**

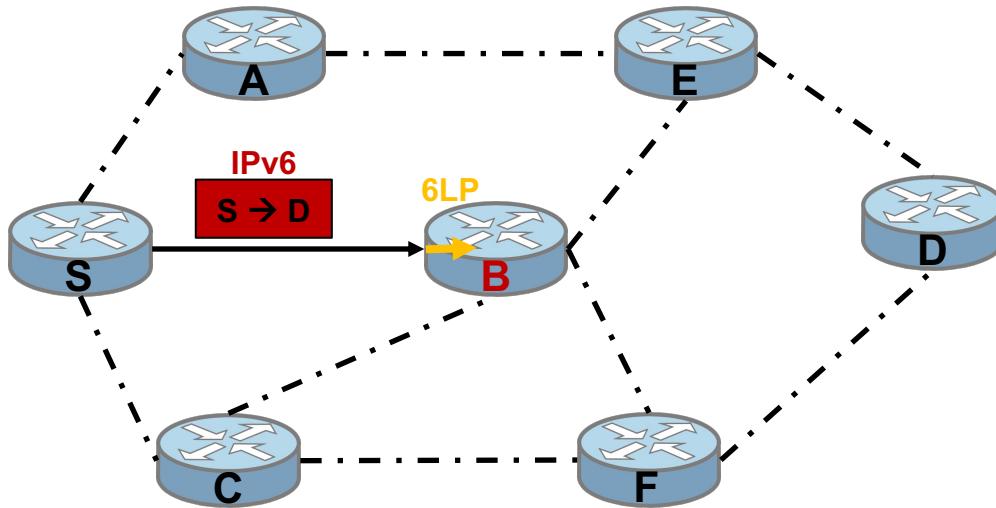
**IPv6: IPv6 Layer**

## Route-Over: Operation Example

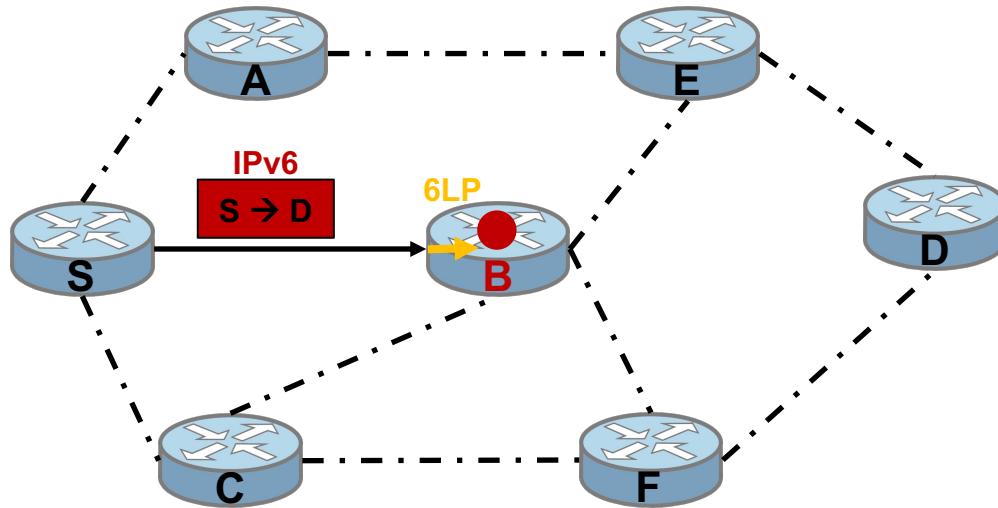


1. Upon the reception of the 1<sup>st</sup> fragment:
  - node B stores the received fragments.

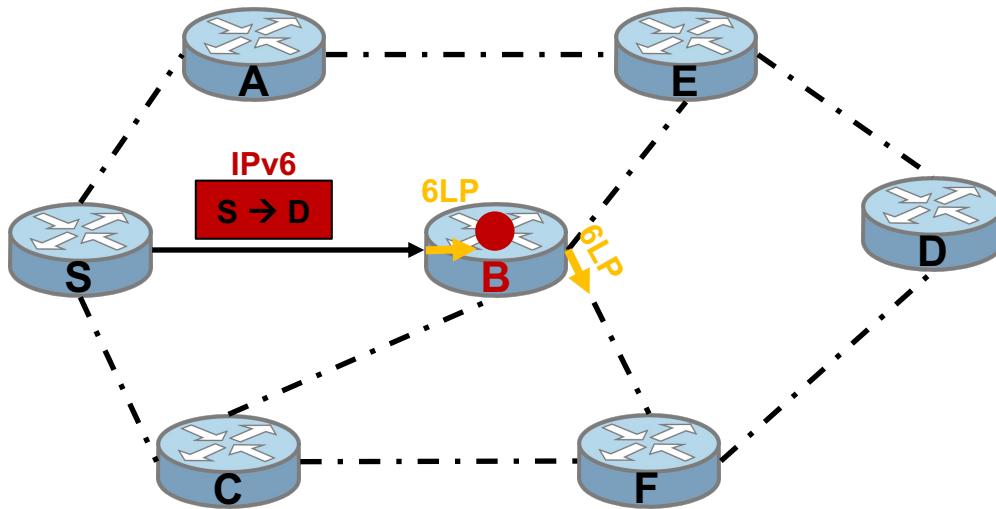
## Route-Over: Operation Example



1. Upon the reception of the 1<sup>st</sup> fragment:
  - node B stores the received fragments.
  - then, it initiates the Reassembly and Decompression operations.

**Route-Over: Operation Example**

2. Then, it checks the IPv6 original source and final destination headers.

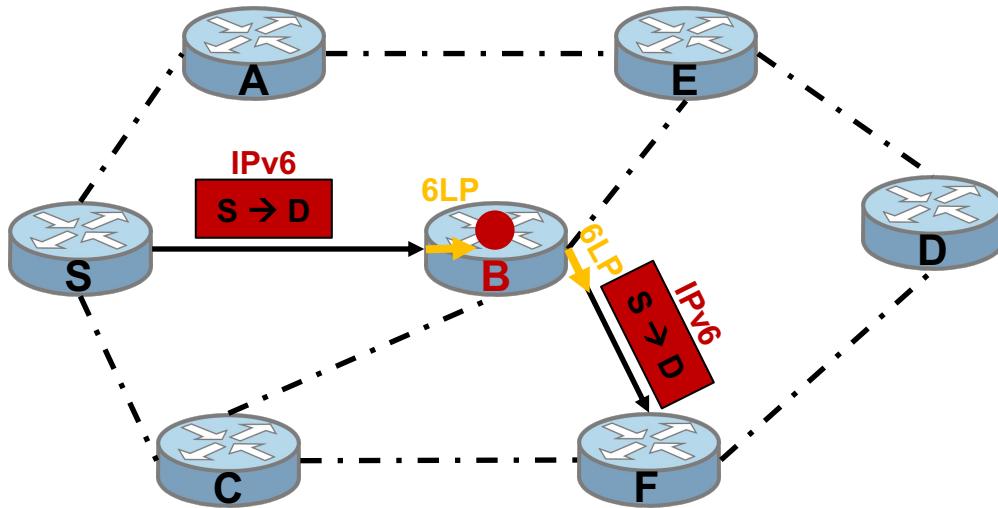
**Route-Over: Operation Example**

3. If node **B** is not the *final destination*, it pushes to the 6LoWPAN layer:
  - for Compression and Fragmentation operations.

**IPv6:** IPv6 Layer

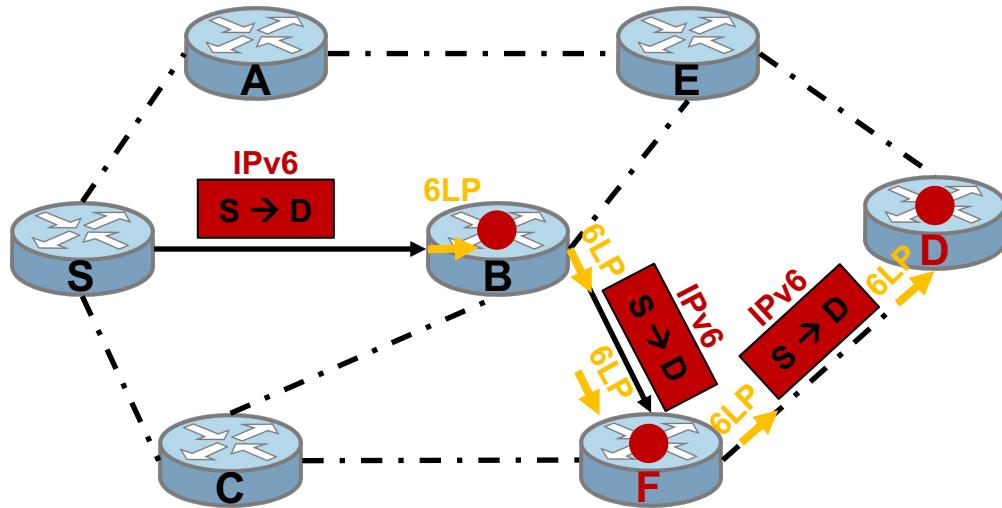
**6LP:** 6LoWPAN Adaptation Layer

## Route-Over: Operation Example



4. Based on its L3 routing table:
  - it will forward the packet to node F.

## Route-Over: Operation Example



**IPv6:** IPv6 Layer

**6LP:** 6LoWPAN Adaptation Layer

# Chapter 3

## Route Over (Per-hop Fragmentation & Reassembly): Issues

## Route-Over: Issues

For more details, please check the following **articles**:

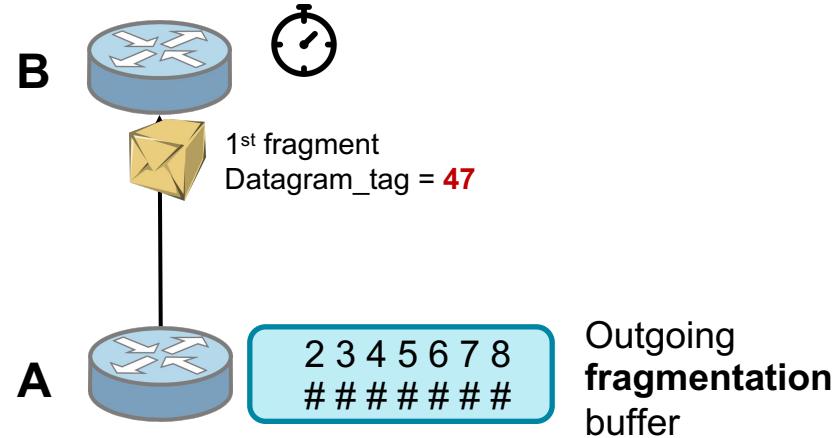
1. G. Z. Papadopoulos, R. Jadhav, P. Thubert and N. Montavont,  
*"Updates on RFC 4944: Fragment Forwarding and Recovery,"*  
*In IEEE Communications Standards Magazine, vol. 3, pp. 54-59, June 2019.*
2. G. Z. Papadopoulos, P. Thubert, S. Tsakalidis and N. Montavont,  
*"RFC 4944: Per-hop Fragmentation and Reassembly Issues"*  
*In Proc. IEEE CSCN 2018 - Paris, France, October 2018*

## Route-Over: Issues (Reliability)

RFC 4944

### Reliability:

- When a node receives the 1<sup>st</sup> fragment, it initiates the **reassembly timer** (i.e., 60 seconds).



## Route-Over: Issues (Reliability)

RFC 4944

### Reliability:

- ▶ When a node receives the 1<sup>st</sup> fragment, it initiates the ***reassembly timer*** (i.e., 60 seconds).
- ▶ If **all fragments are not received** during this *reassembly period*:
  - the **reassembly** of the IPv6 packet **is not possible**.

**Route-Over: Issues (Latency)**

RFC 4944

## Latency:

- ▶ The reassembly at each hop increases the delay:
  - *each device is required to “wait” for 60 seconds.*

	Source	Forwarder 1	Forwarder 2	Destination
T = 0	F F F			
T = 1	F F (F)	F		
T = 2	F (F)	FF		
T = 3	(F)	FFF		
T = 4		FF (F)	F	
T = 5		F (F)	FF	
T = 6		(F)	FFF	
T = 7			FF (F)	F
T = 8			F (F)	FF
T = 9			(F)	FFF

## Route-Over: Issues (Latency)

RFC 4944

### Latency:

- ▶ The reassembly at each hop increases the delay:
  - *each device is required to “wait” for 60 seconds.*
- ▶ There is an additional computation time for the fragmentation operation.

	Source	Forwarder 1	Forwarder 2	Destination
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T = 1	F F (F)	F		
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T = 5		F (F)	FF	
T = 6		(F)	FFF	
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T = 8			F (F)	FF
T = 9			(F)	FFF

## Route-Over: Issues (Latency)

RFC 4944

### Latency:

- ▶ The reassembly at each hop increases the delay:
  - *each device is required to “wait” for 60 seconds.*
- ▶ There is an additional computation time for the fragmentation operation.
- ▶ Thus, **the more hops** in the path toward the destination is the **higher** the end-to-end **delay**.

	Source	Forwarder 1	Forwarder 2	Destination
T = 0	F F F			
T = 1	F F (F)	F		
T = 2	F (F)	FF		
T = 3	(F)	FFF		
T = 4		FF (F)	F	
T = 5		F (F)	FF	
T = 6		(F)	FFF	
T = 7			FF (F)	F
T = 8			F (F)	FF
T = 9			(F)	FFF

## Route-Over: Issues (Resource Usage)

### Resource Usage:

- ▶ The fragmentation process requires **large** incoming reassembly **buffers** (1280 bytes or more).

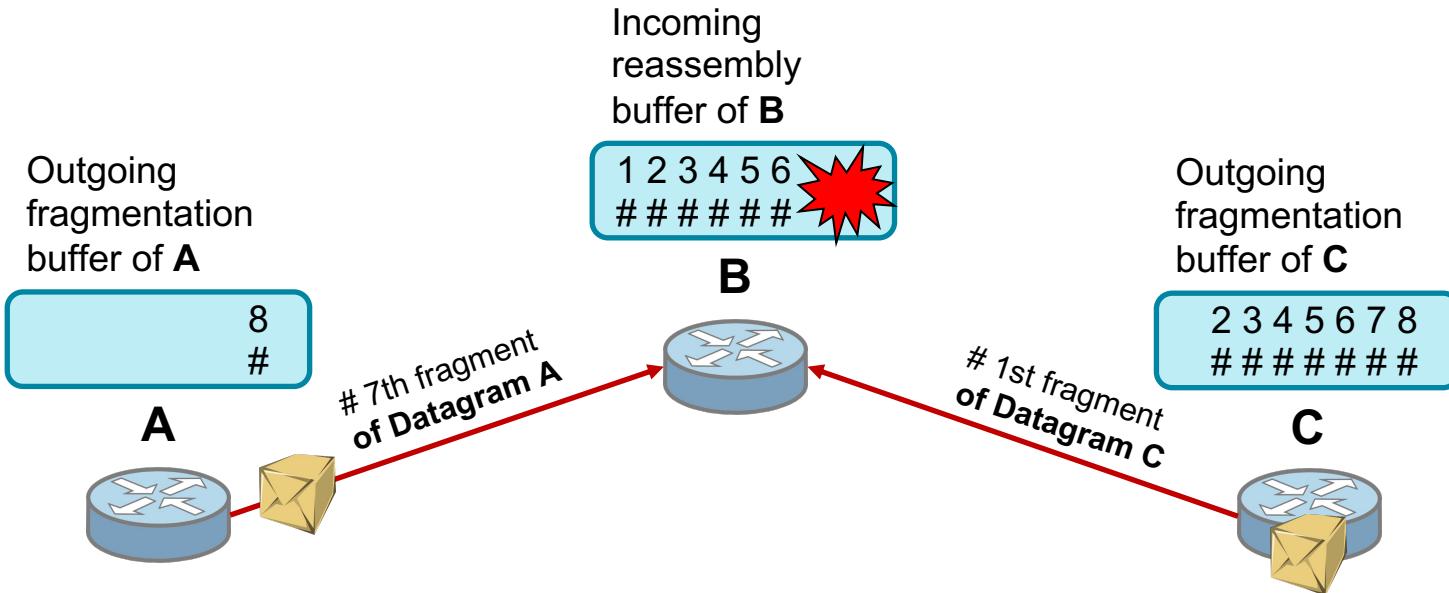
RFC 4944

## Route-Over: Issues (Resource Usage)

RFC 4944

### Resource Usage:

- ▶ The fragmentation process requires **large** incoming reassembly **buffers** (1280 bytes or more).
- ▶ It will be possible to reassemble only very few complete IPv6 packets.

**Route-Over: Issues (Resource Usage)**

## Route-Over: Issues (Resource Usage)

RFC 4944

### Resource Usage:

- ▶ The fragmentation process requires **large** incoming reassembly **buffers** (1280 bytes or more).
- ▶ It will be possible to reassemble only very few complete IPv6 packets.
- ▶ Thus, such issues introduces more losses in a multi-hop network.

**Route-Over: Issues (Implementation or *datagram\_tag*)****Implementation:**

- ▶ Only the **1<sup>st</sup> fragment** of an IPv6 packet contains the source and destination **IPv6 addresses**:

RFC 4944

**Route-Over: Issues (*Implementation or datagram\_tag*)**

RFC 4944

### Implementation:

- ▶ Only the **1<sup>st</sup> fragment** of an IPv6 packet contains the source and destination **IPv6 addresses**:
- ▶ The subsequent fragments are routed based on the *datagram\_tag*.

**Route-Over: Issues (Implementation or *datagram\_tag*)**

RFC 4944

**Implementation:**

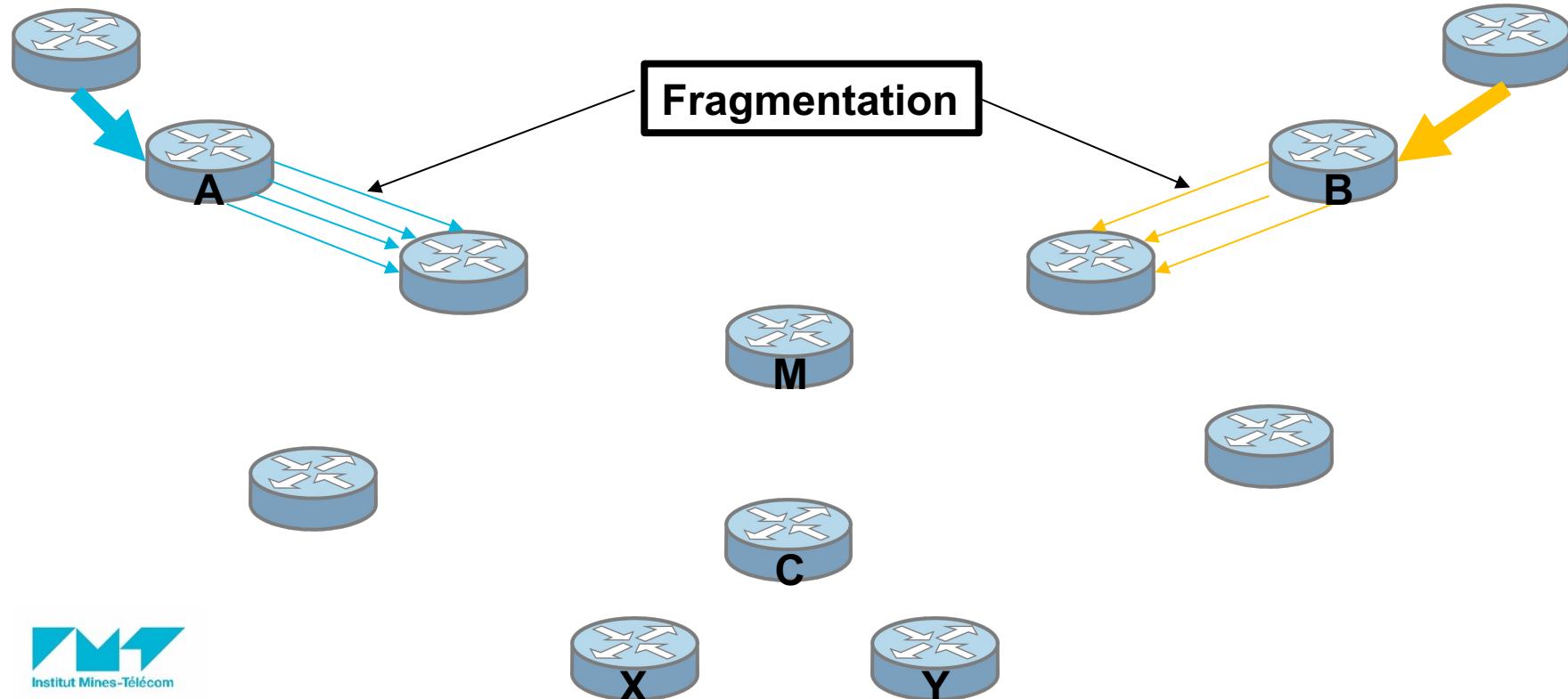
- ▶ Only the **1<sup>st</sup> fragment** of an IPv6 packet contains the source and destination **IPv6 addresses**:
- ▶ The subsequent fragments are routed based on the ***datagram\_tag***.
- ▶ A ***datagram\_tag* is unique only** to the 6LoWPAN original **source** and final **destination nodes**.

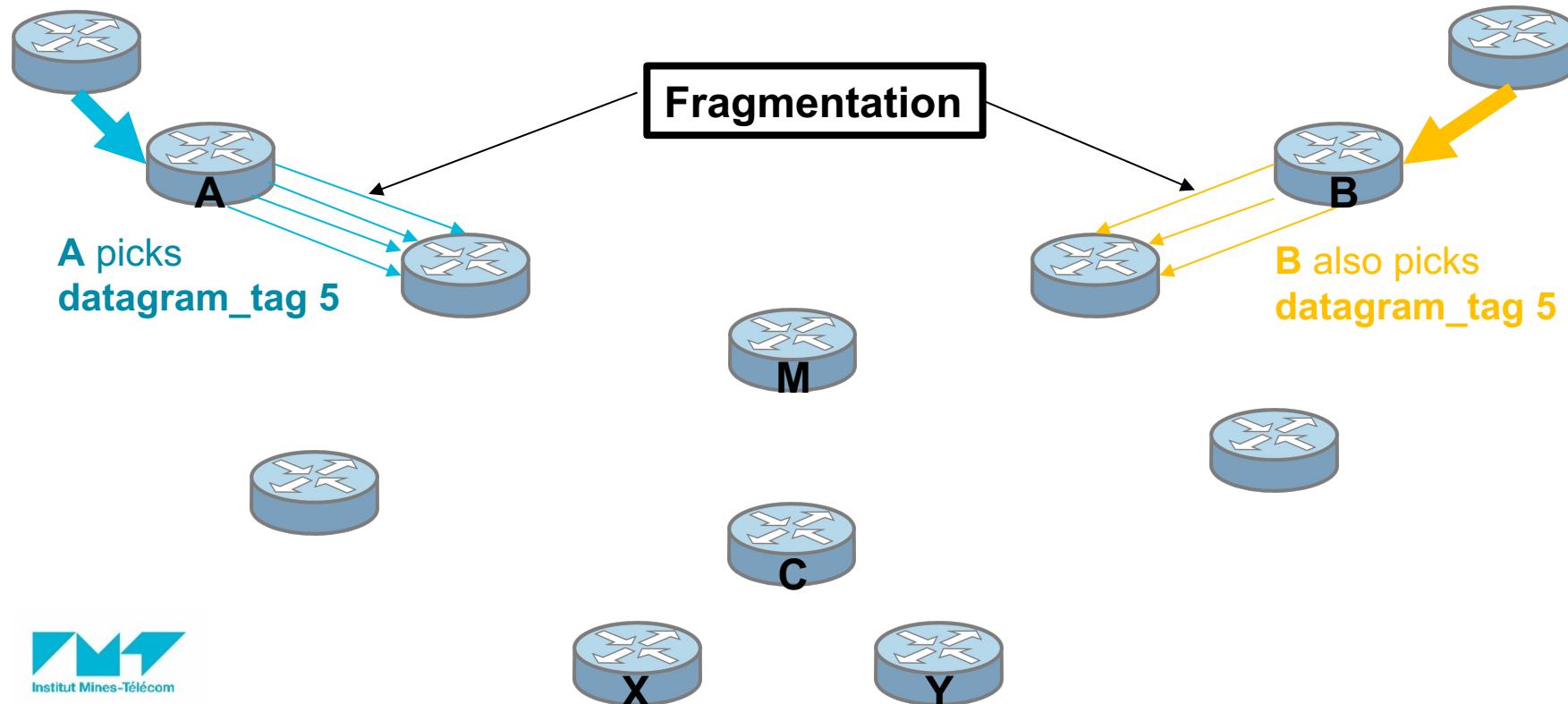
**Route-Over: Issues (Implementation or *datagram\_tag*)**

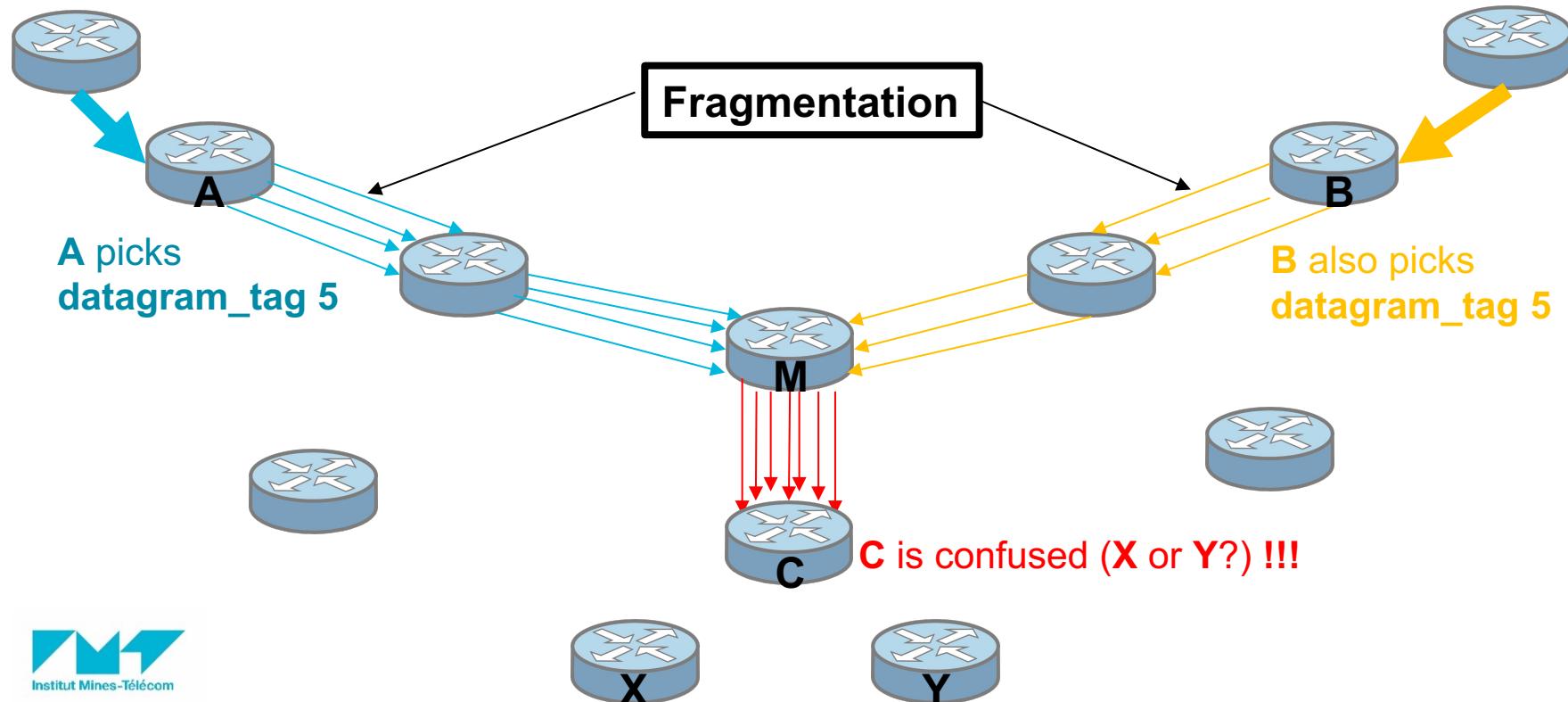
RFC 4944

### Implementation:

- ▶ Only the **1<sup>st</sup> fragment** of an IPv6 packet contains the source and destination **IPv6 addresses**:
- ▶ The subsequent fragments are routed based on the ***datagram\_tag***.
- ▶ A ***datagram\_tag* is unique only** to the 6LoWPAN original **source** and final **destination nodes**.
- ▶ Thus, two different traffic flows may be tagged with the same value.

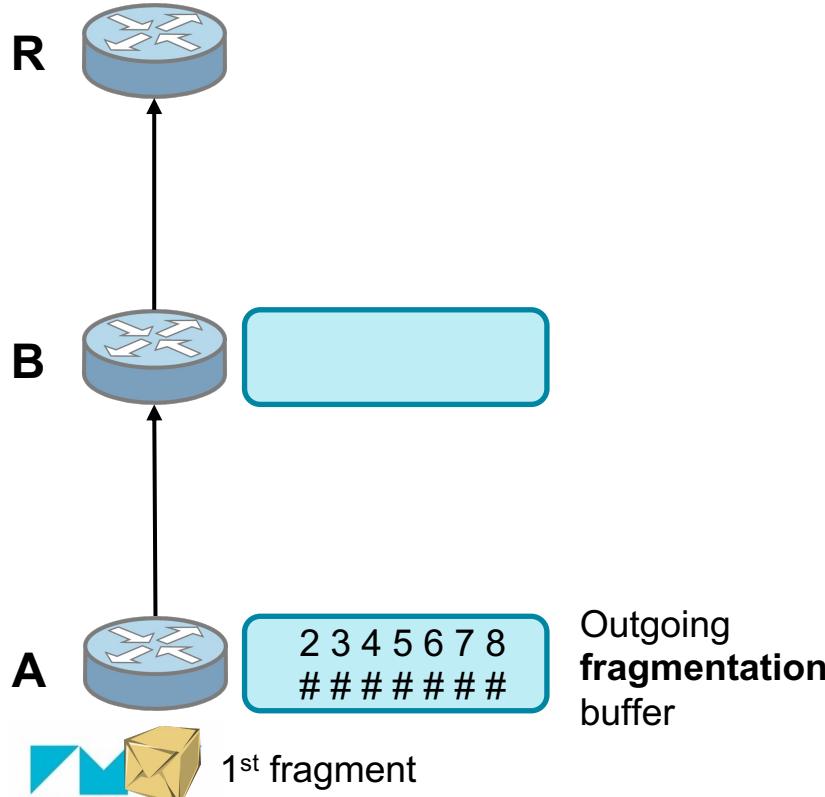
Route-Over: Issues (Implementation or *datagram\_tag*)

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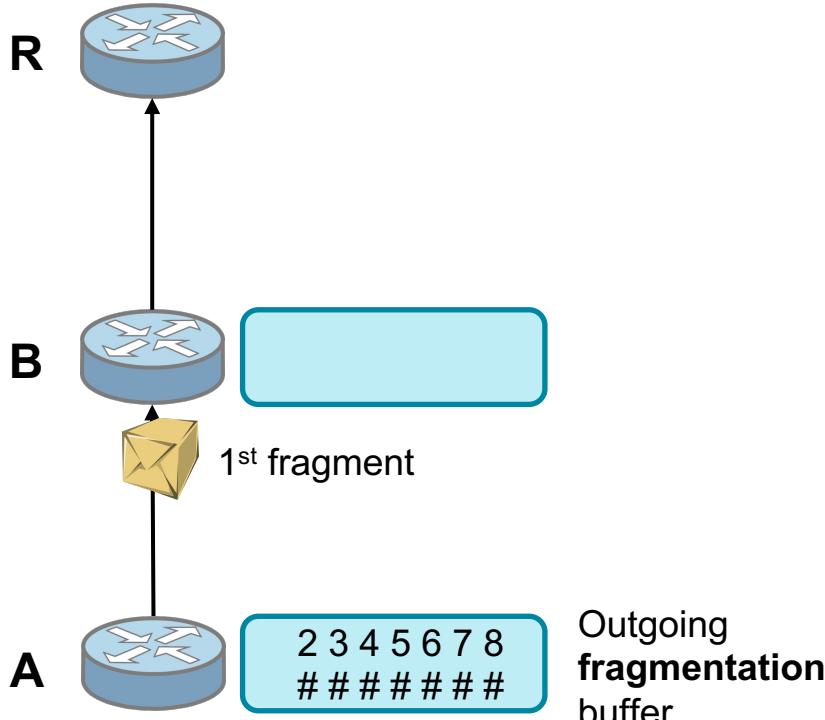
# Chapter 4

## RFC 8930: 6LoWPAN Fragment Forwarding (6LFF)



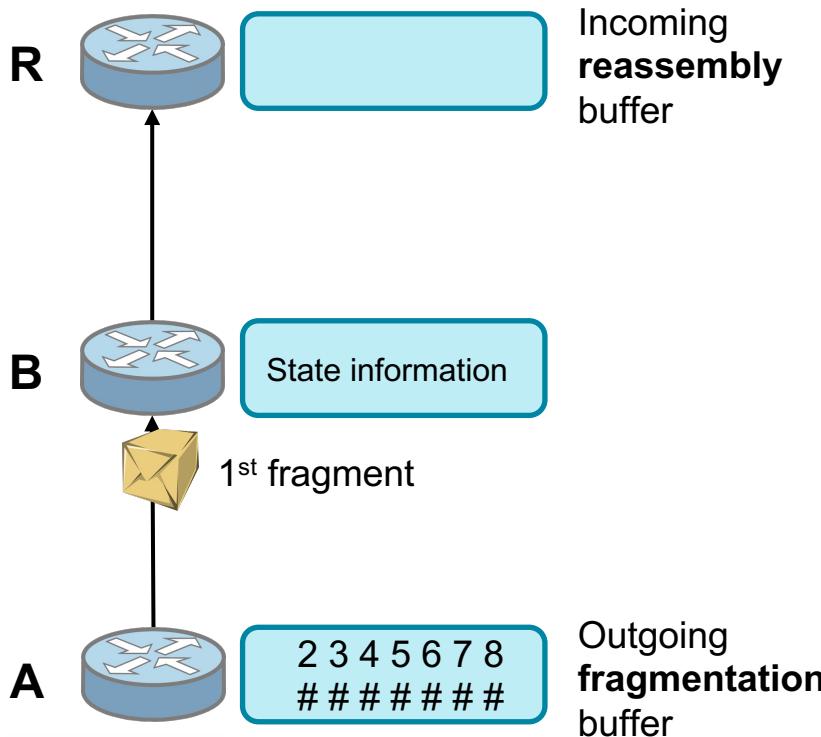
## 6LFF at a glance:

- ▶ The routing decision is made at the 1<sup>st</sup> fragment.



## 6LFF at a glance:

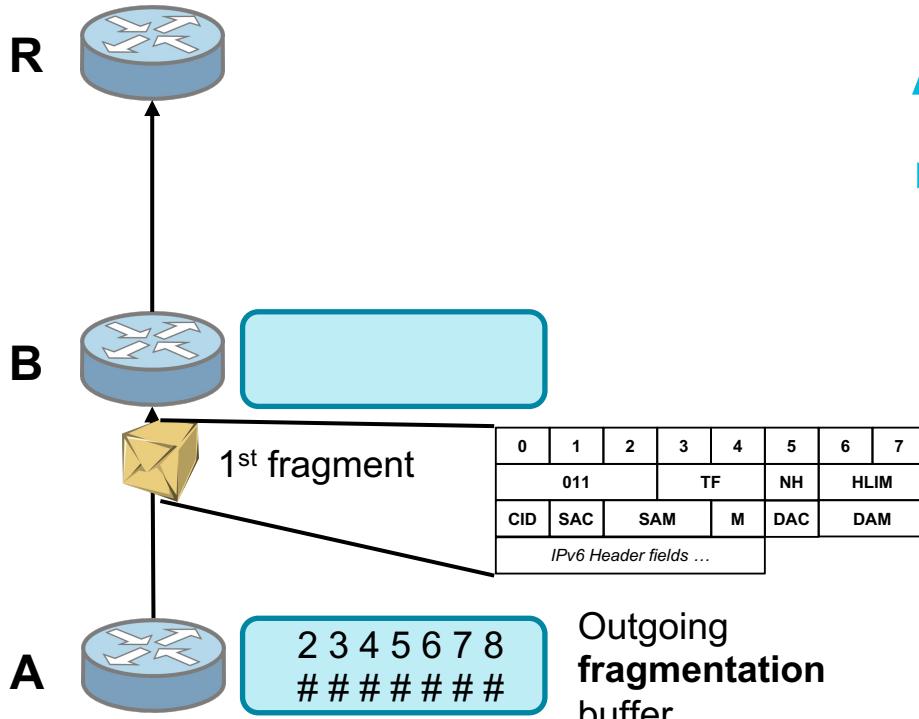
- ▶ The routing decision is made at the 1<sup>st</sup> fragment.
- ▶ The 1<sup>st</sup> fragment carries the ***IPv6 compressed header*** ***destination address***.



## 6LFF at a glance:

- ▶ The routing decision is made at the 1<sup>st</sup> fragment.
- ▶ The 1<sup>st</sup> fragment carries the ***IPv6 compressed header*** ***destination address***.
- ▶ The received fragments are forwarded immediately!

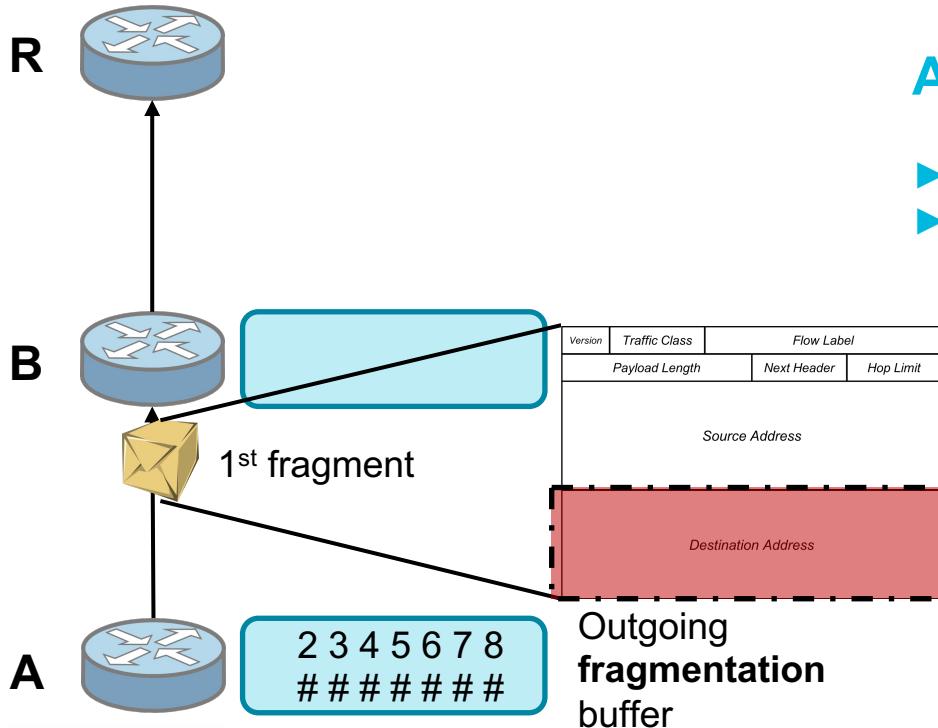
## The Operation of the 6LFF Mechanism



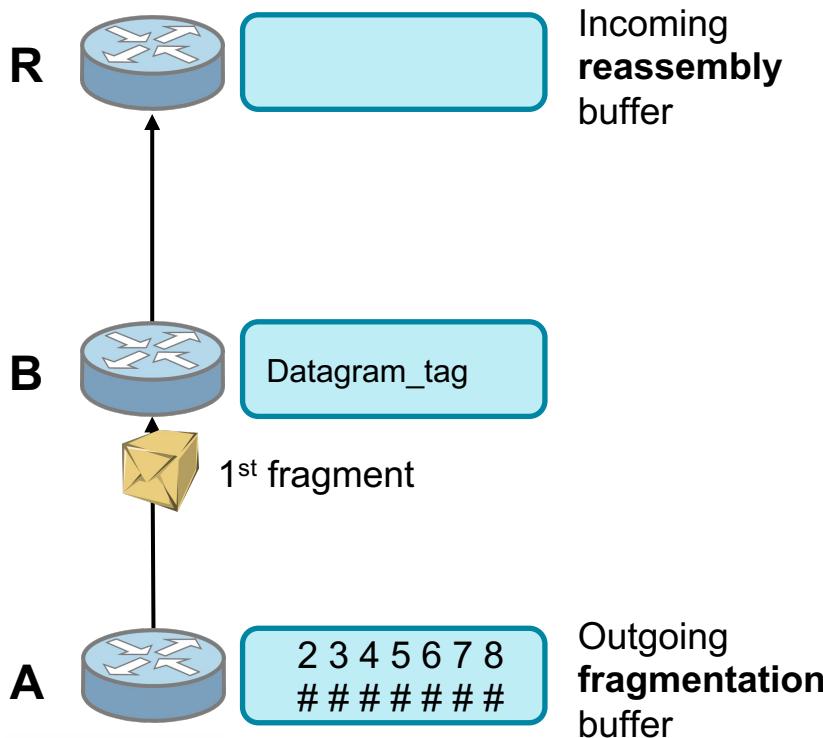
**At the reception of the 1<sup>st</sup> fragment:**

- Decompress the ***IPv6 compressed header***.

## The Operation of the 6LFF Mechanism



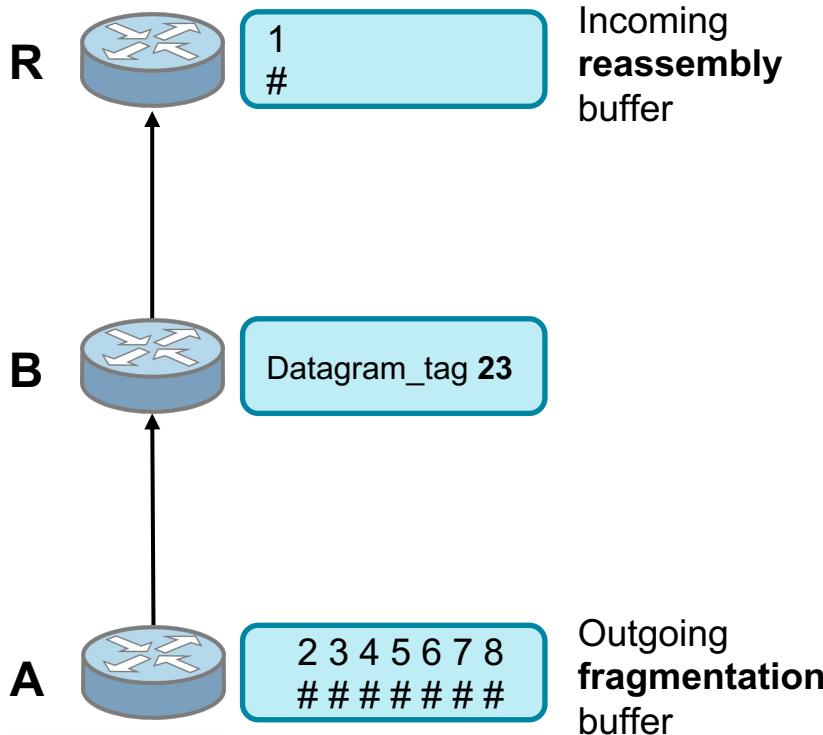
## The Operation of the 6LFF Mechanism



**At the reception of the 1<sup>st</sup> fragment:**

- ▶ Decompress the ***IPv6 compressed header***.
- ▶ To identify the IPv6 address of the next hop.
- ▶ Forward the fragment to the next hop.

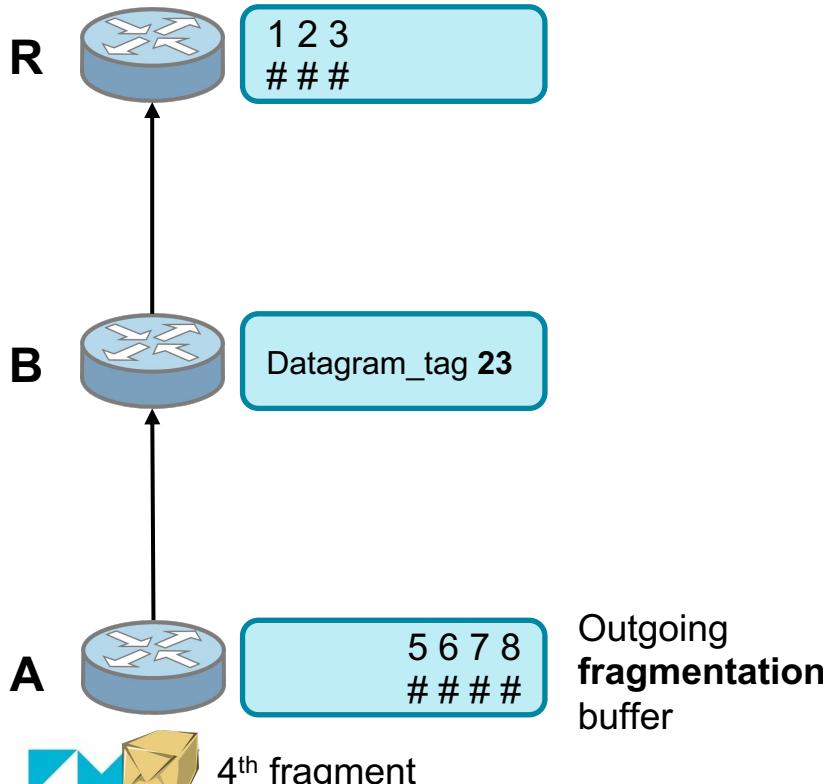
## The Operation of the 6LFF Mechanism



**At the reception of the 1<sup>st</sup> fragment:**

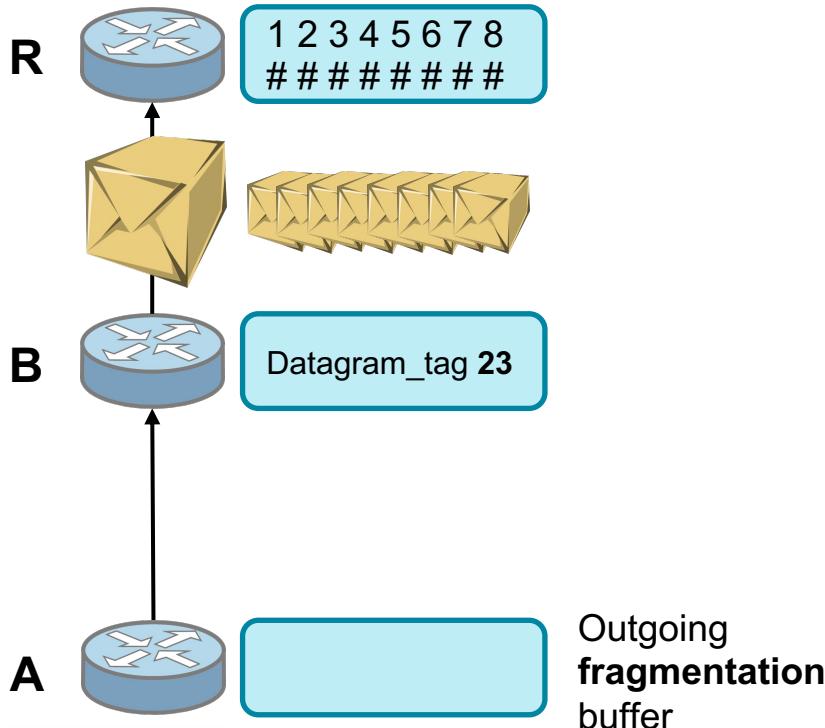
- ▶ Decompress the ***IPv6 compressed header***.
- ▶ To identify the IPv6 address of the next hop.
- ▶ Forward the fragment to the next hop.
- ▶ Register the ***datagram\_tag*** of the fragment.

## The Operation of the 6LFF Mechanism



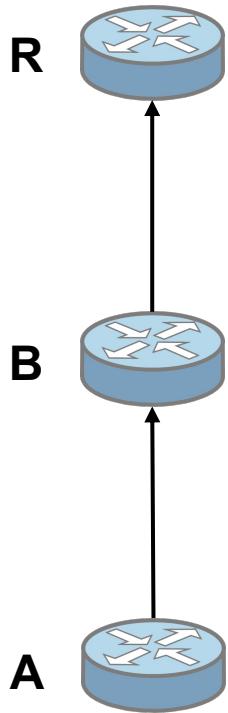
**At the reception of the subsequent fragments:**

## The Operation of the 6LFF Mechanism



**At the reception of the subsequent fragments:**

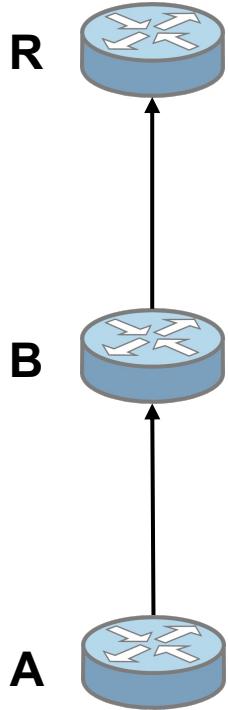
- ▶ The IPv6 packet is reassembled when **all** the fragments have arrived at the destination.

**Virtual Reassembly Buffer (VRB)**

Node B's VRB table.			
incoming	outgoing		

**VRB at a glance:**

- ▶ The operation is similar to ***switching table***.

**Virtual Reassembly Buffer (VRB)**

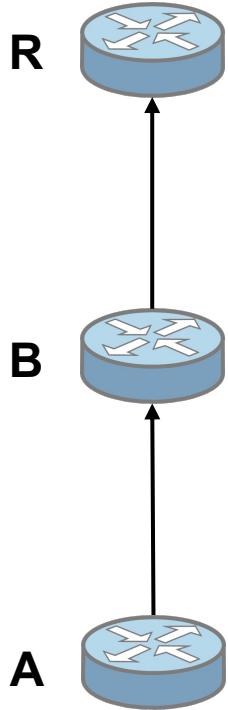
Node B's VRB table.

incoming	outgoing
entry 1 (IPv6 packet 1)	
entry 2 (IPv6 packet 2)	
entry ... (IPv6 packet ...)	
entry n (IPv6 packet n)	

**VRB at a glance:**

- ▶ The operation is similar to **switching table**.
- ▶ The entries correspond to IPv6 packets.

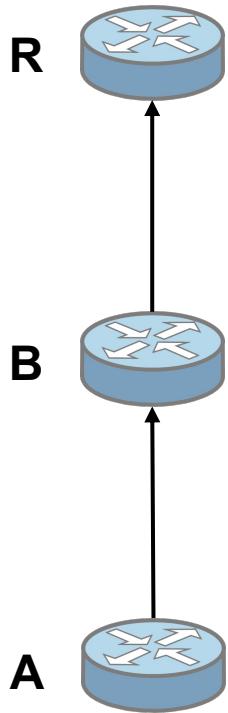
## Virtual Reassembly Buffer (VRB)



Node B's VRB table.	
incoming	outgoing
entry 1 (empty)	
entry 2 (empty)	
entry ... (empty)	
entry $n$ (empty)	

### VRB at a glance:

- ▶ The operation is similar to ***switching table***.
- ▶ The entries correspond to IPv6 packets.
- ▶ The VRBs have a maximum pre-allocated memory.

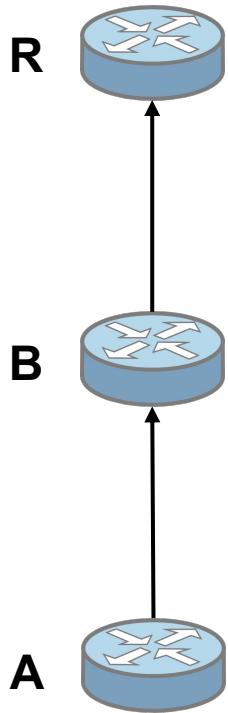
**Virtual Reassembly Buffer (VRB)**

Node B's VRB table.

incoming	outgoing		

**A VRB entry:**

- ▶ Each VRB entry is a tuple of 4 elements:

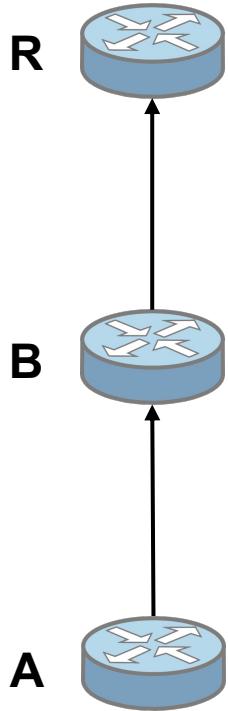
**Virtual Reassembly Buffer (VRB)**

Node B's VRB table.

incoming	outgoing		
L2 src			

**A VRB entry:**

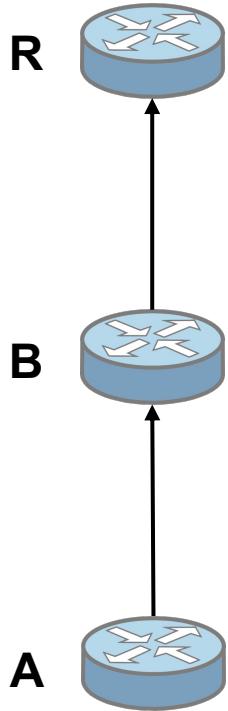
- ▶ Each VRB entry is a tuple of 4 elements:
  - link-layer address of the previous hop.

**Virtual Reassembly Buffer (VRB)**

Node B's VRB table.			
incoming	outgoing		
L2 src	tag		

**A VRB entry:**

- ▶ Each VRB entry is a tuple of 4 elements:
  - link-layer address of the previous hop.
  - locally unique datagram\_tag of the incoming fragment.

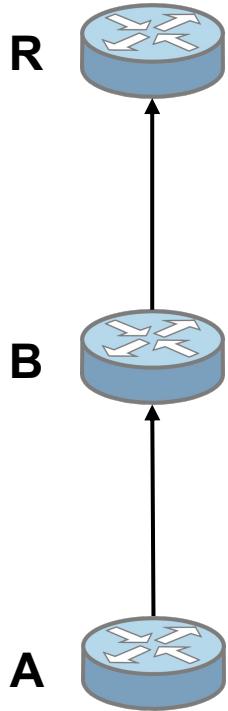
**Virtual Reassembly Buffer (VRB)**

Node B's VRB table.

incoming		outgoing	
L2 src	tag	L2 dest	

**A VRB entry:**

- ▶ Each VRB entry is a tuple of 4 elements:
  - link-layer address of the previous hop.
  - locally unique datagram\_tag of the incoming fragment.
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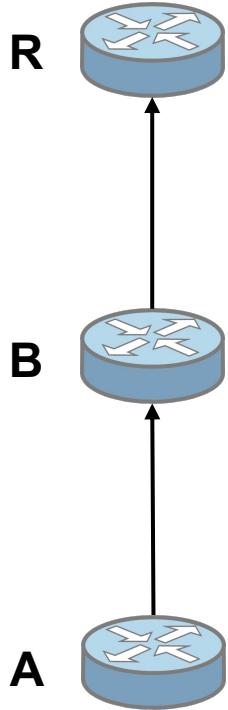
**Virtual Reassembly Buffer (VRB)**

Node B's VRB table.

incoming		outgoing	
L2 src	tag	L2 dest	tag

**A VRB entry:**

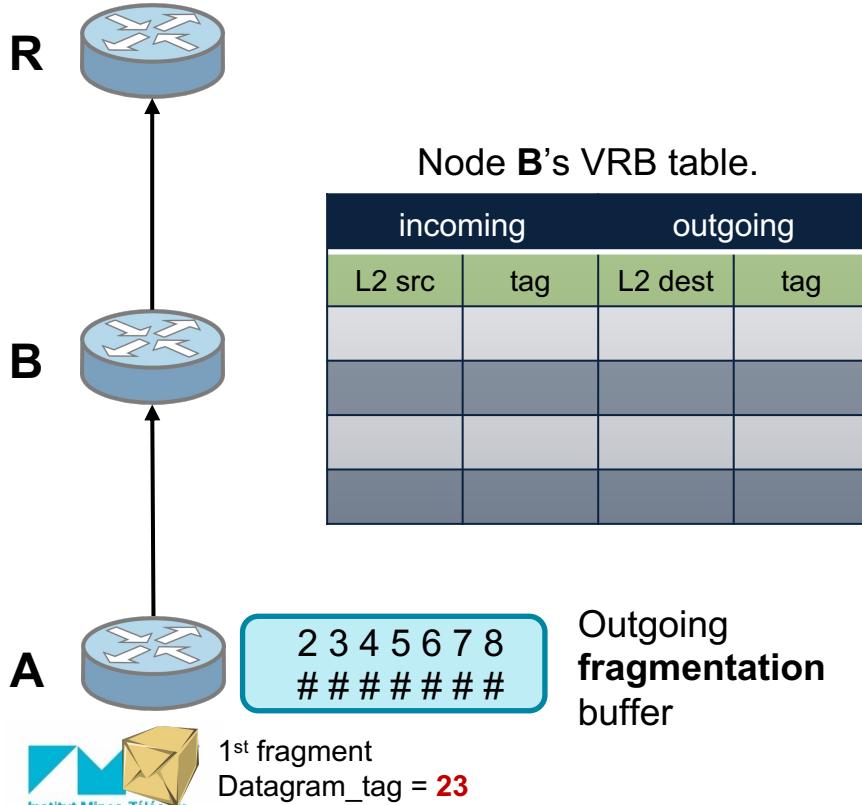
- ▶ Each VRB entry is a tuple of 4 elements:
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  - locally unique datagram\_tag of the incoming fragment.
  - link-layer address of the next hop.
  - locally unique datagram\_tag for the outgoing fragment.

**Virtual Reassembly Buffer (VRB)**

incoming		outgoing	
L2 src	tag	L2 dest	tag

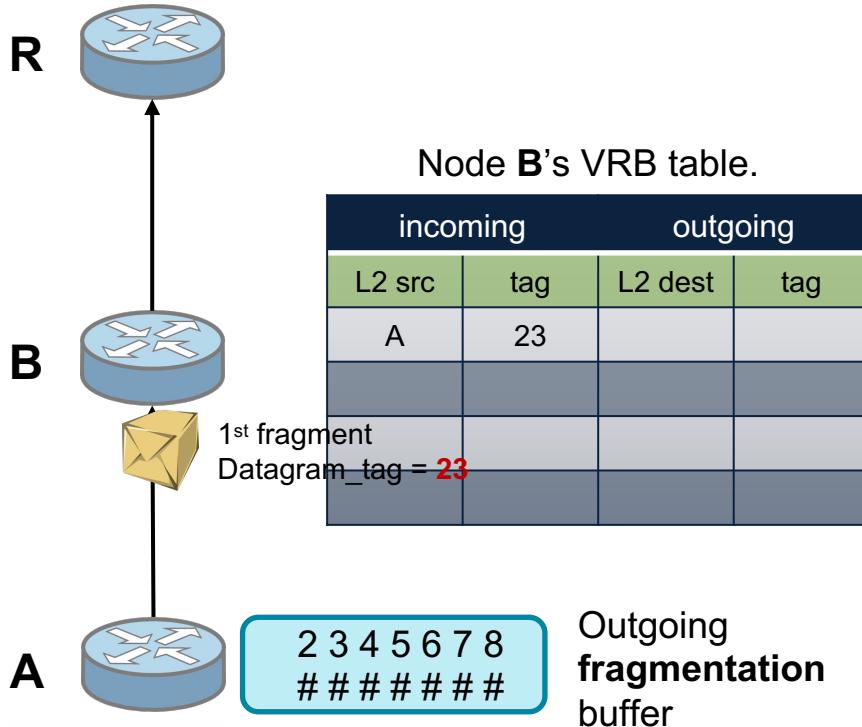
**A VRB entry:**

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  - link-layer address of the previous hop.
  - locally unique datagram\_tag of the incoming fragment.
  - link-layer address of the next hop.
  - locally unique datagram\_tag for the outgoing fragment.
- ▶ Each VRB entry requires 20 bytes.

**VRB Operation**

At the reception of the 1<sup>st</sup> fragment:

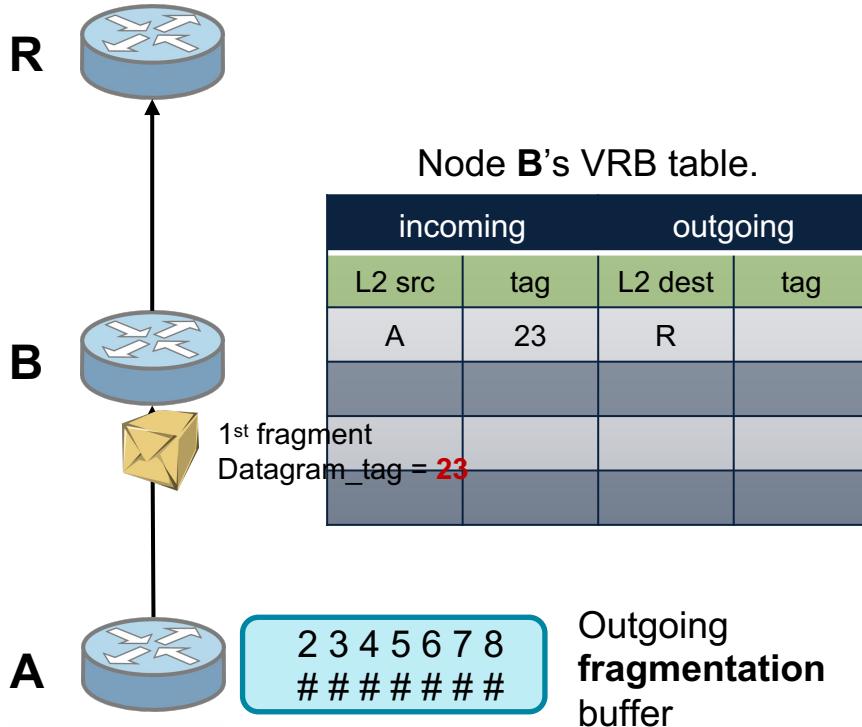
## VRB Operation



**At the reception of the 1<sup>st</sup> fragment:**

- ▶ Record the *link-layer address* and *datagram\_tag* of the **previous hop**.

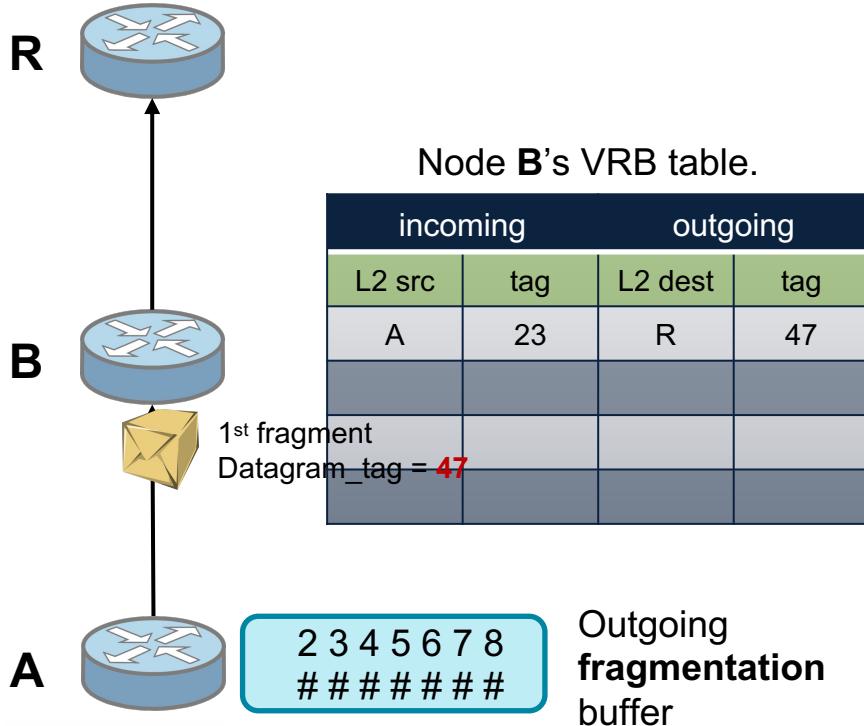
## VRB Operation



**At the reception of the 1<sup>st</sup> fragment:**

- ▶ Record the *link-layer address* and *datagram\_tag* of the **previous hop**.
- ▶ Determine the *link-layer address* of the **next hop**.

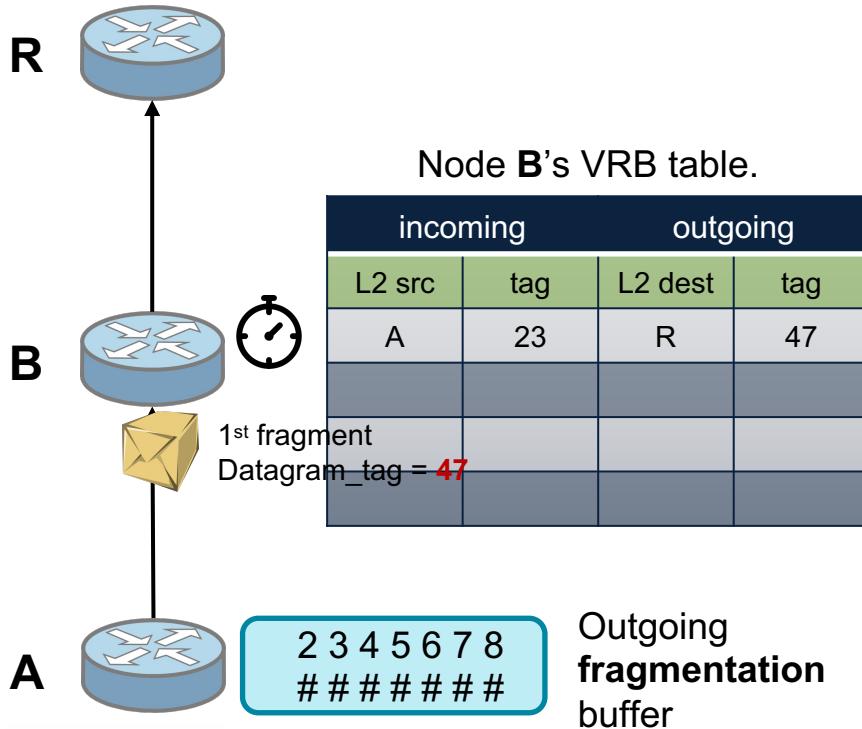
## VRB Operation



**At the reception of the 1<sup>st</sup> fragment:**

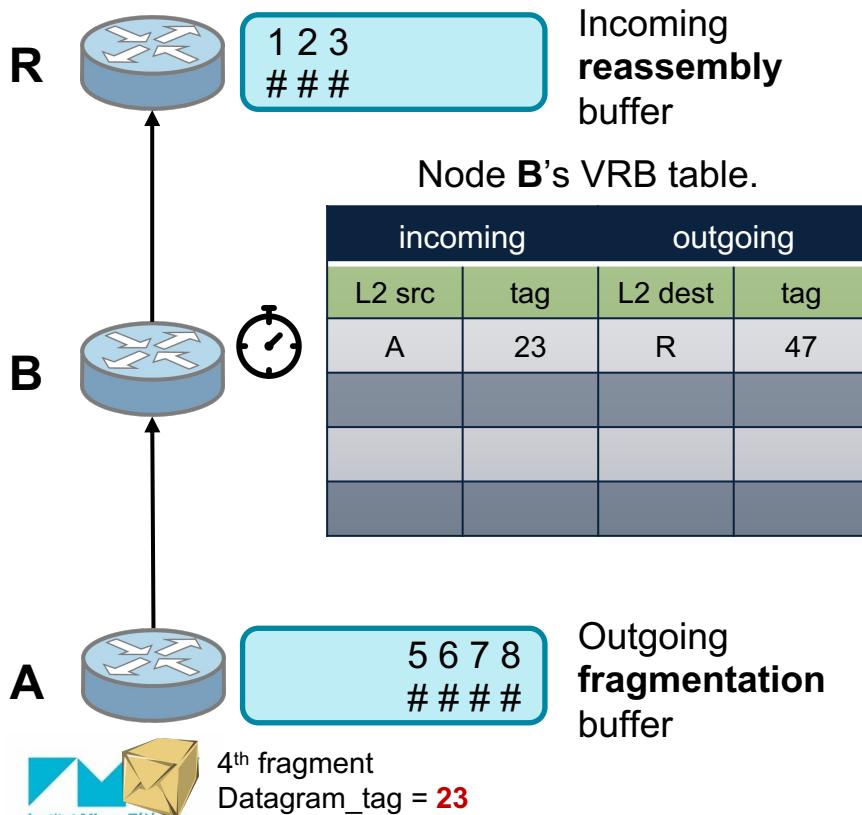
- ▶ Record the *link-layer address* and *datagram\_tag* of the **previous hop**.
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- ▶ Pick a **new unique datagram\_tag** for the **next hop**.

## VRB Operation

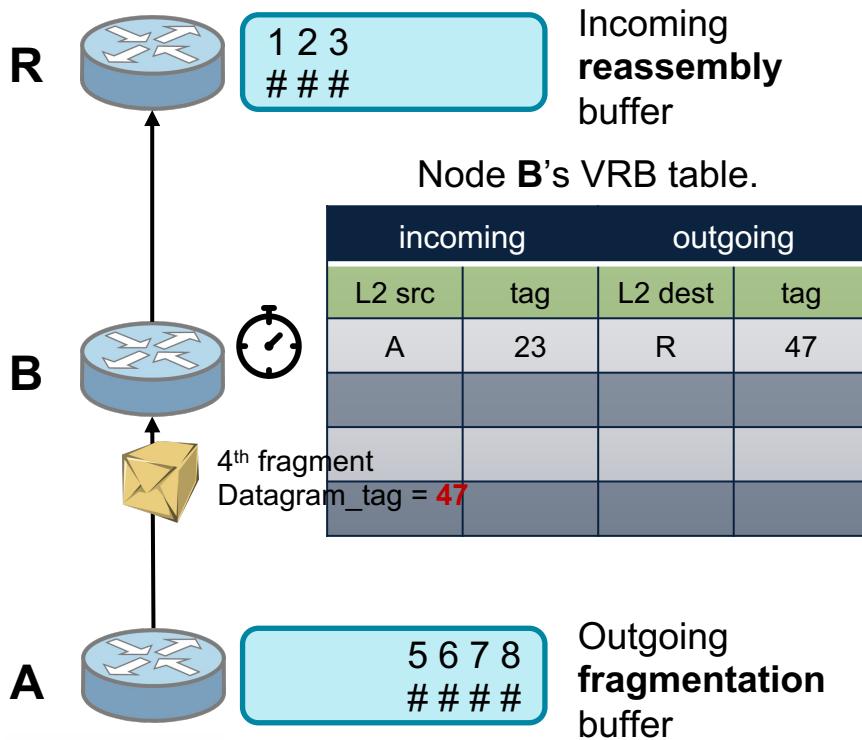


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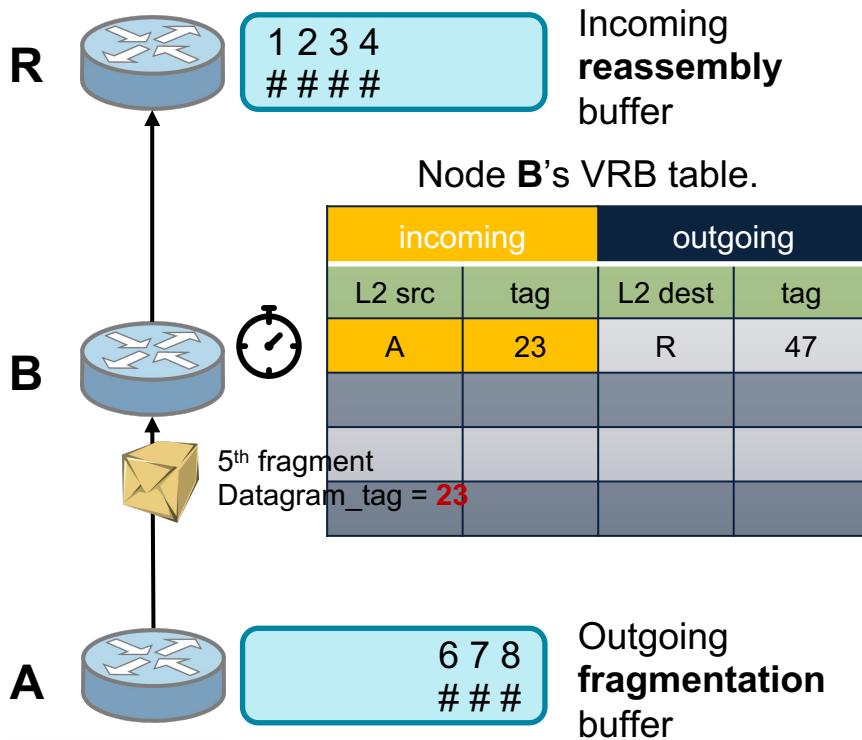
- ▶ Record the *link-layer address* and *datagram\_tag* of the **previous hop**.
- ▶ Determine the *link-layer address* of the **next hop**.
- ▶ Pick a **new unique datagram\_tag** for the **next hop**.
- ▶ Set a timer to allow discarding a **partially reassembled packet** after some timeout.

**VRB Operation**

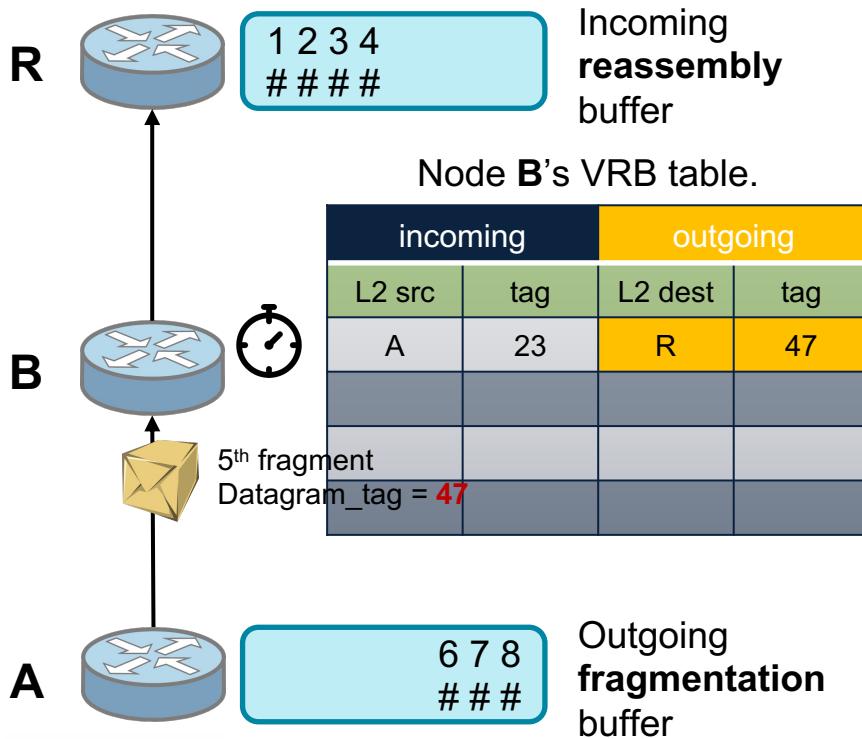
RFC 8930

**VRB Operation**

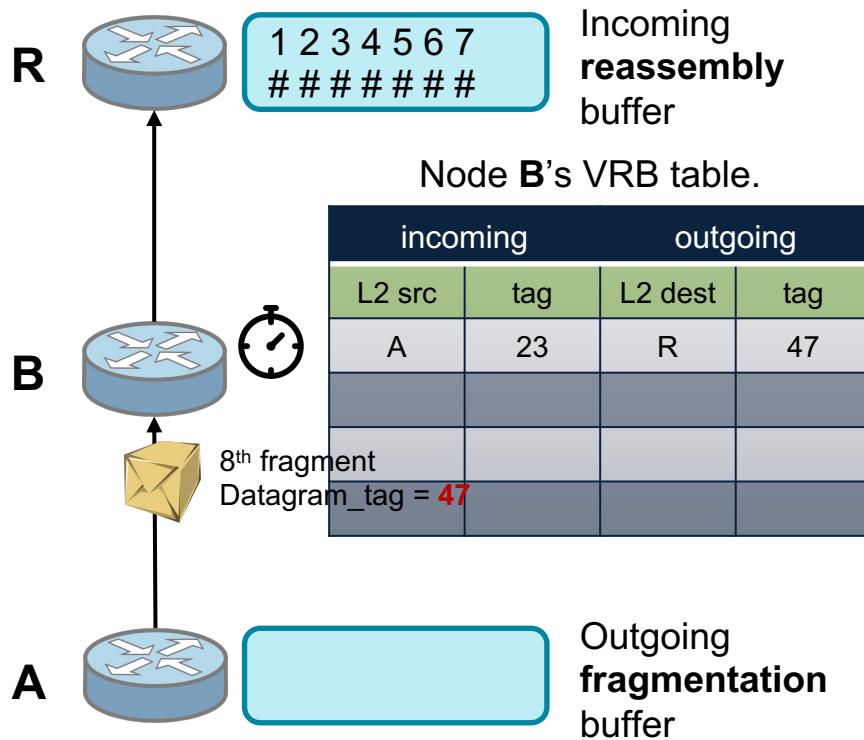
RFC 8930

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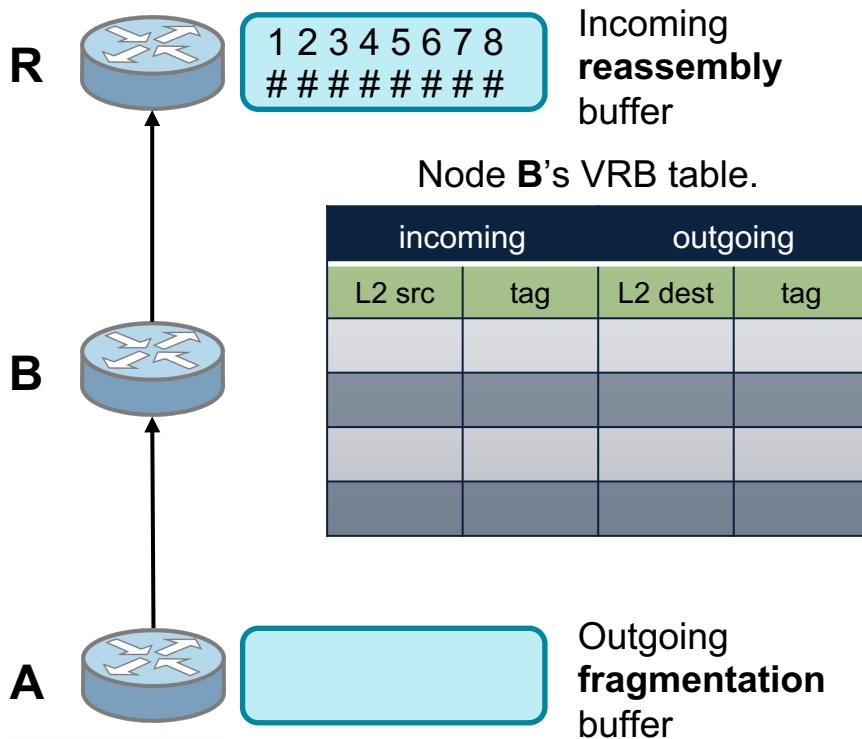
RFC 8930

**VRB Operation**

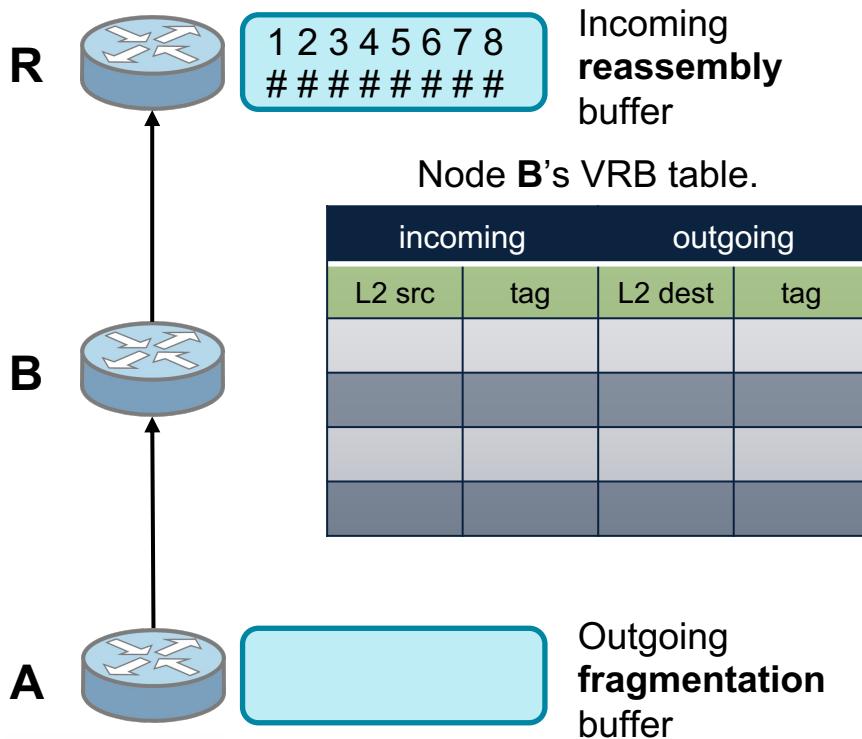
RFC 8930

**VRB Operation**

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RFC 8930

## Advantages

- ▶ The **end-to-end latency** should be greatly reduced.

RFC 8930

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RFC 8930

- ▶ The **end-to-end latency** should be greatly reduced.
- ▶ The **end-to-end network reliability** should be improved.
- ▶ The **datagram\_tag issue** is solved.

## Drawbacks

- ▶ There is **no Fragment Recovery** built-in.

RFC 8930

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RFC 8930

- ▶ There is **no Fragment Recovery** built-in.
  - It introduces unnecessary traffic in the network.
- ▶ **6LFF does not support per-fragment routing.**
- ▶ The **size of the VRB table necessarily remains finite.**

RFC 8930



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# 6LoWPAN: Fragmentation & Reassembly, Frame Delivery Modes, & Fragment Forwarding

Georgios Z. PAPADOPOULOS, Associate Professor, IMT Atlantique

[georgios.papadopoulos@imt-atlantique.fr](mailto:georgios.papadopoulos@imt-atlantique.fr)

[www.georgiospapadopoulos.com](http://www.georgiospapadopoulos.com)

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**About this Course**  
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By presenting the building blocks of the IoT network architecture, this MOOC will help learners adapt to the fast changing communications and networking environment of IoT.

The IoT world represents billions of sophisticated objects, such as sensors, actuators and meters, that are deployed nearly everywhere, in homes, hospitals, factories, cities, and are connected to the Internet. However, they come with limited capacity in terms of memory storage.

**SHOW ALL**

**WHAT YOU WILL LEARN**

- ✓ How to schedule a collision free communication between two devices (with TSCH protocol, MSF)
- ✓ How to compress, fragment and reassemble IPv6 data packets adapted to IoT constraints (with 6Lowpan, 6LFF)
- ✓ How make connected devices learn their best path toward a given destination (with RPL protocol)

**SKILLS YOU WILL GAIN**

Routing Protocol Internet Of Things (IOT) Wireless Network Architecture

**Flexible deadlines**  
Reset deadlines in accordance to your schedule.

**Shareable Certificate**  
Earn a Certificate upon completion

**100% online**  
Start instantly and learn at your own schedule.

**Intermediate Level**  
- Basic notions of programming for the Laboratory sessions (i.e., C, Python).

**Approx. 16 hours to complete**

**English**  
Subtitles: English

**■ Outline of the MOOC:**

- Week 1: Welcome & MAC Methods (i.e., TSCH)
- Week 2: 6TiSCH
- Week 3: IPv6 & 6LoWPAN
- Week 4: RPL

**■ Educational Team:**

- Georgios Z. Papadopoulos
- Nicolas Montavont
- Géraldine Texier
- Remous-Aris Koutsiamanis

**■ Interviews from the industrial community:**

- Pascal Thubert, Cisco Systems
- Thomas Watteyne, Analog Devices, Falco
- Rémi Dubaele, ENEDIS